

BGP Techniques for Internet Service Providers

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Presentation Slides

Will be available on

ftp://ftp-eng.cisco.com

/pfs/seminars/MENOG2-BGP-Techniques.pdf

And on the MENOG2 website

Feel free to ask questions any time

BGP Techniques for Internet Service Providers

- BGP Basics
- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



BGP Basics

What is BGP?

Border Gateway Protocol

 A Routing Protocol used to exchange routing information between different networks

Exterior gateway protocol

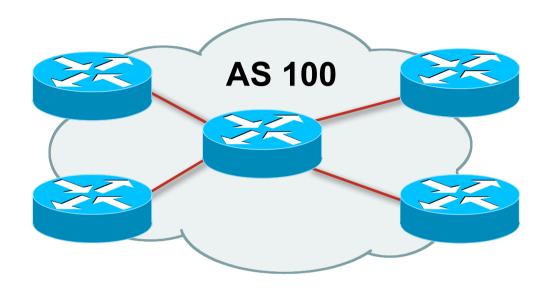
Described in RFC4271

RFC4276 gives an implementation report on BGP RFC4277 describes operational experiences using BGP

The Autonomous System is BGP's fundamental operating unit

It is used to uniquely identify networks with a common routing policy

Autonomous System (AS)



- Collection of networks with same routing policy
- Single routing protocol
- Usually under single ownership, trust and administrative control
- Identified by a unique number (ASN)

Autonomous System Number (ASN)

An ASN is a 16 bit integer

1-64511 are for use on the public Internet

64512-65534 are for private use only

0 and 65535 are reserved

ASNs are now extended to 32 bit!

RFC4893 is standards document describing 32-bit ASNs

Representation still under discussion:

32-bit notation or "16.16" notation

Latter documented in Internet Draft:

draft-michaelson-4byte-as-representation-04.txt

AS 23456 is used to represent 32-bit ASNs in 16-bit ASN world

Autonomous System Number (ASN)

 ASNs are distributed by the Regional Internet Registries

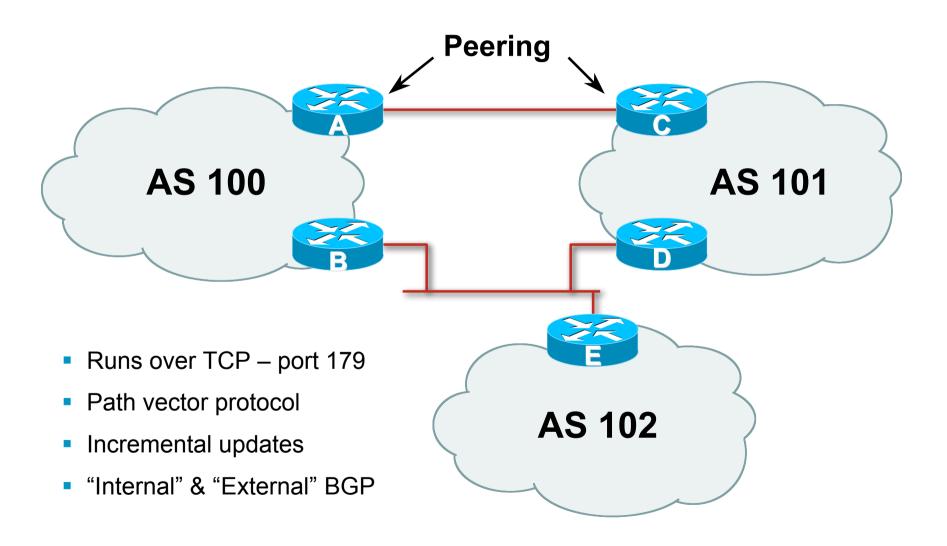
They are also available from upstream ISPs who are members of one of the RIRs

 Current 16-bit ASN allocations up to 45055 have been made to the RIRs

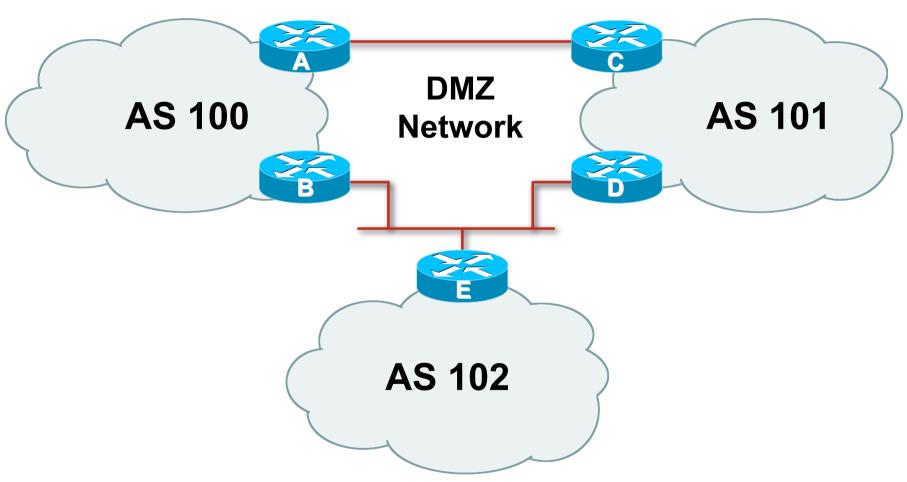
Around 26600 are visible on the Internet

- The RIRs also have received 1024 32-bit ASNs each
 5 are visible on the Internet (early adopters)
- See www.iana.org/assignments/as-numbers

BGP Basics



Demarcation Zone (DMZ)



Shared network between ASes

BGP General Operation

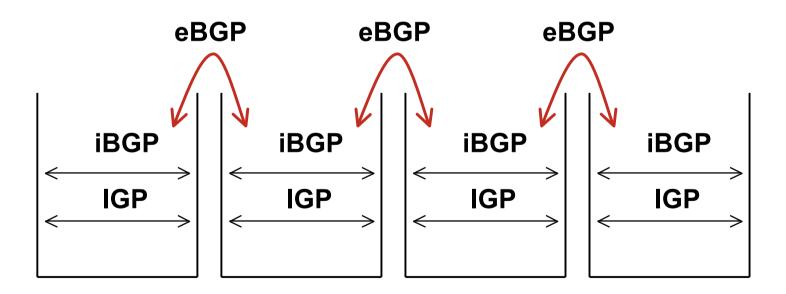
- Learns multiple paths via internal and external BGP speakers
- Picks the best path and installs in the forwarding table
- Best path is sent to external BGP neighbours
- Policies are applied by influencing the best path selection

eBGP & iBGP

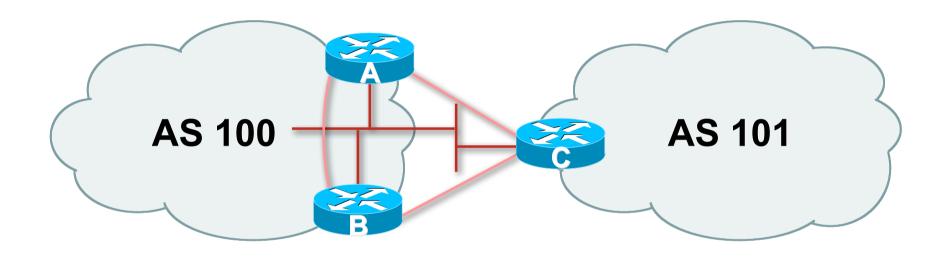
- BGP used internally (iBGP) and externally (eBGP)
- iBGP used to carry some/all Internet prefixes across ISP backbone ISP's customer prefixes
- eBGP used to exchange prefixes with other ASes implement routing policy

BGP/IGP model used in ISP networks

Model representation



External BGP Peering (eBGP)

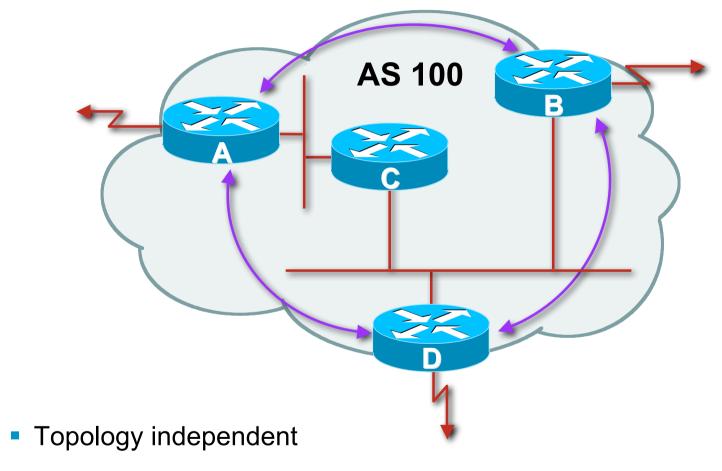


- Between BGP speakers in different AS
- Should be directly connected
- Never run an IGP between eBGP peers

Internal BGP (iBGP)

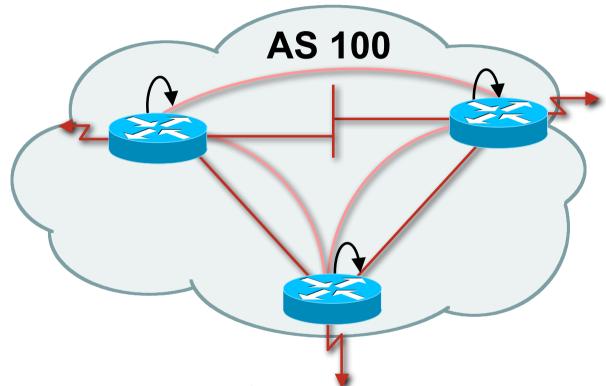
- BGP peer within the same AS
- Not required to be directly connected
 IGP takes care of inter-BGP speaker connectivity
- iBGP speakers must to be fully meshed:
 - They originate connected networks
 - They pass on prefixes learned from outside the ASN
 - They do not pass on prefixes learned from other iBGP speakers

Internal BGP Peering (iBGP)



 Each iBGP speaker must peer with every other iBGP speaker in the AS

Peering to Loopback Interfaces



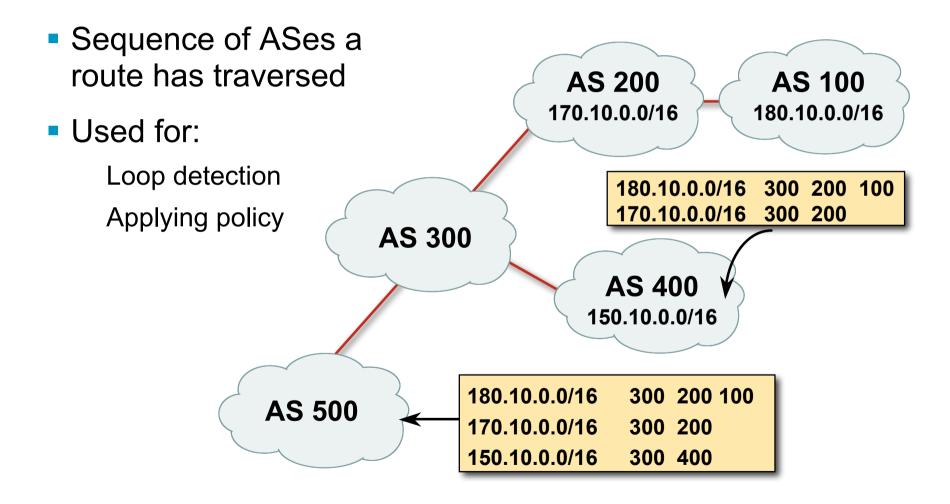
- Peer with loop-back interface
 Loop-back interface does not go down ever!
- Do not want iBGP session to depend on state of a single interface or the physical topology



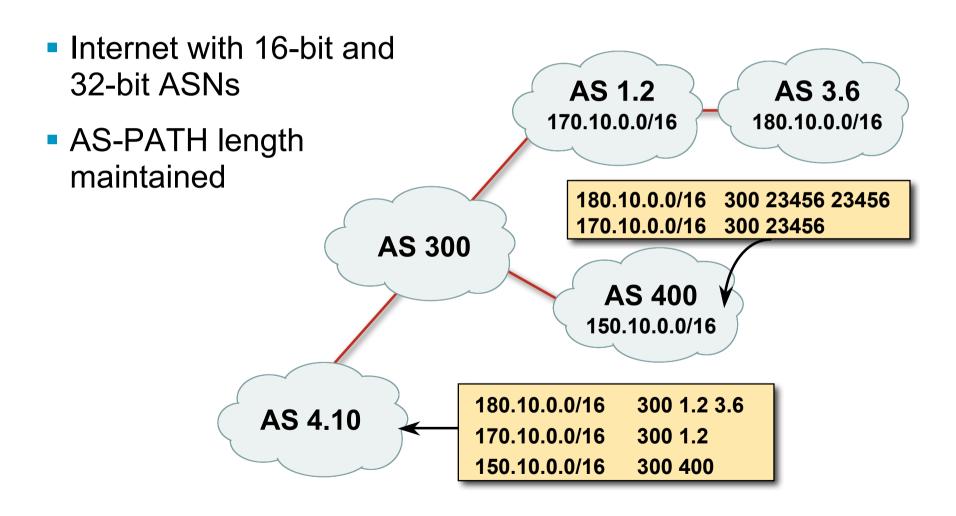
BGP Attributes

Information about BGP

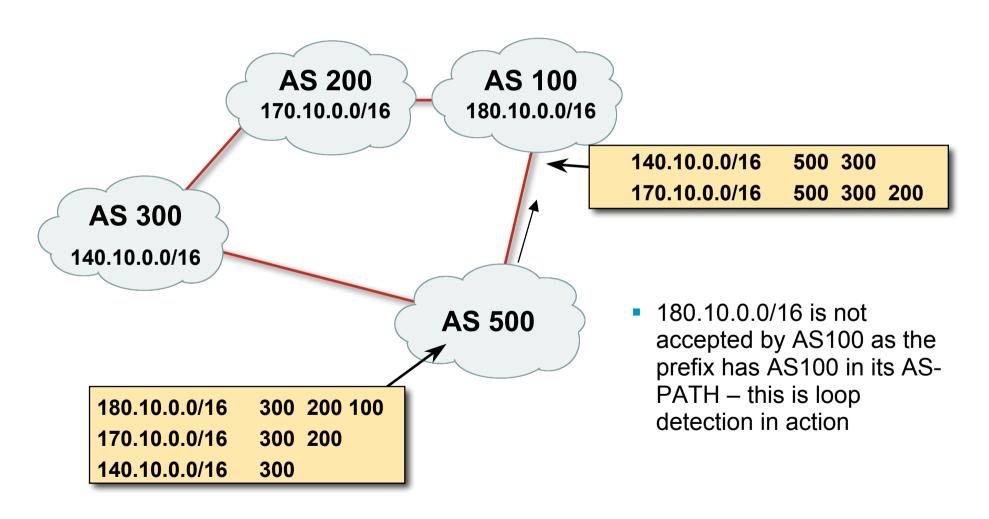
AS-Path



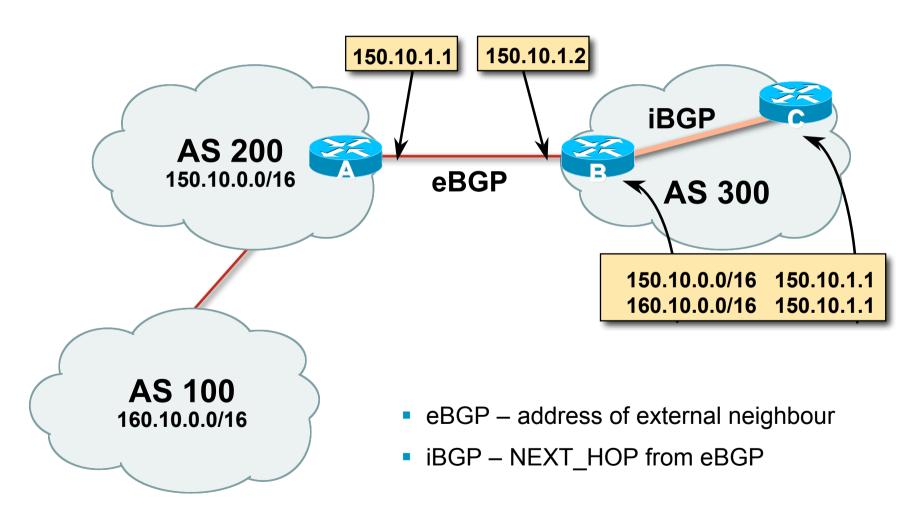
AS-Path (with 16 and 32-bit ASNs)



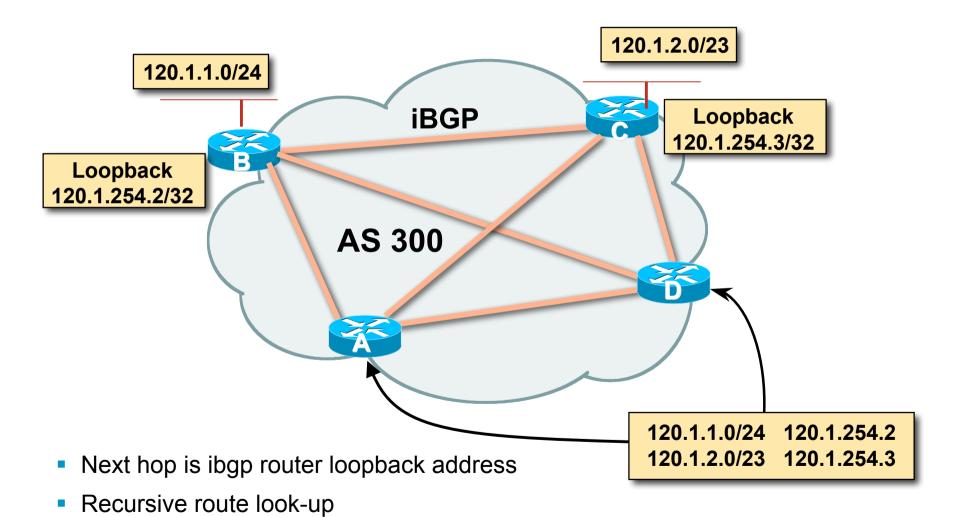
AS-Path loop detection



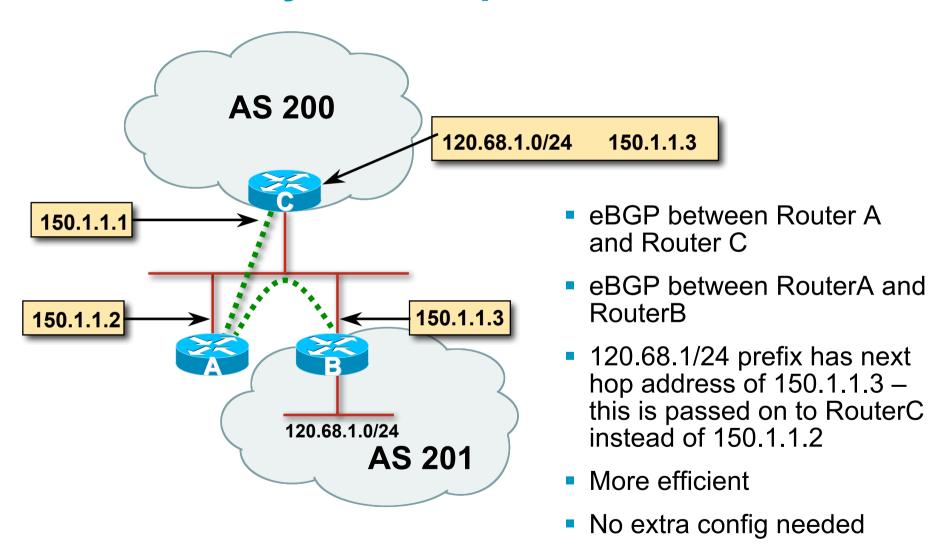
Next Hop



iBGP Next Hop



Third Party Next Hop



Next Hop Best Practice

 BGP default is for external next-hop to be propagated unchanged to iBGP peers

This means that IGP has to carry external next-hops

Forgetting means external network is invisible

With many eBGP peers, it is unnecessary extra load on IGP

 ISP Best Practice is to change external next-hop to be that of the local router

Next Hop (Summary)

- IGP should carry route to next hops
- Recursive route look-up
- Unlinks BGP from actual physical topology
- Change external next hops to that of local router
- Allows IGP to make intelligent forwarding decision

Origin

- Conveys the origin of the prefix
- Historical attribute
 Used in transition from EGP to BGP
- Influences best path selection
- Three values: IGP, EGP, incomplete

IGP – generated by BGP network statement

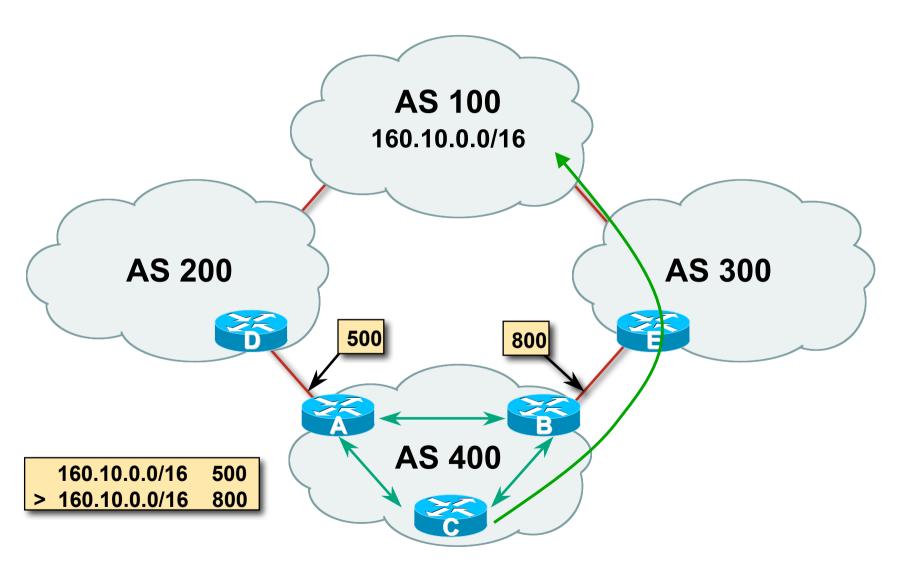
EGP – generated by EGP

incomplete - redistributed from another routing protocol

Aggregator

- Conveys the IP address of the router or BGP speaker generating the aggregate route
- Useful for debugging purposes
- Does not influence best path selection

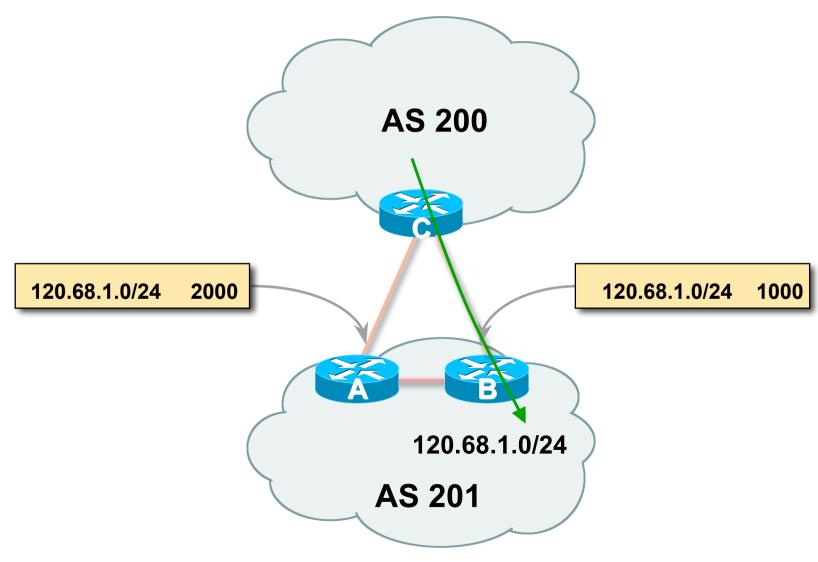
Local Preference



Local Preference

- Local to an AS non-transitive
 Default local preference is 100 (IOS)
- Used to influence BGP path selection determines best path for outbound traffic
- Path with highest local preference wins

Multi-Exit Discriminator (MED)



Multi-Exit Discriminator

- Inter-AS non-transitive & optional attribute
- Used to convey the relative preference of entry points determines best path for inbound traffic
- Comparable if paths are from same AS
 Implementations have a knob to allow comparisons of MEDs from different ASes
- Path with lowest MED wins
- Absence of MED attribute implies MED value of zero (RFC4271)

Multi-Exit Discriminator "metric confusion"

MED is non-transitive and optional attribute

Some implementations send learned MEDs to iBGP peers by default, others do not

Some implementations send MEDs to eBGP peers by default, others do not

 Default metric varies according to vendor implementation

Original BGP spec (RFC1771) made no recommendation

Some implementations said that absence of metric was equivalent to 0

Other implementations said that absence of metric was equivalent to 2³²-1 (highest possible) or 2³²-2

Potential for "metric confusion"

Community

- Communities are described in RFC1997
 Transitive and Optional Attribute
- 32 bit integer

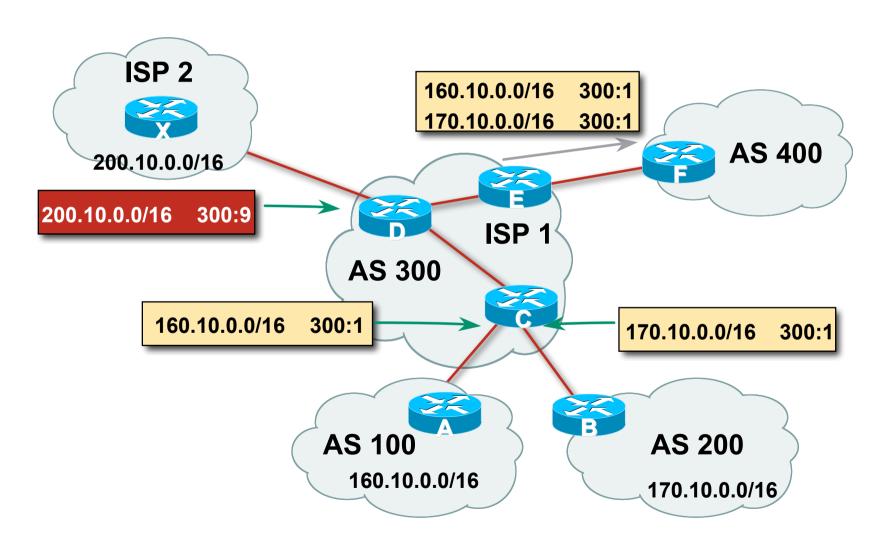
Represented as two 16 bit integers (RFC1998)

Common format is <local-ASN>:xx

0:0 to 0:65535 and 65535:0 to 65535:65535 are reserved

- Used to group destinations
 Each destination could be member of multiple communities
- Very useful in applying policies within and between ASes

Community



Well-Known Communities

Several well known communities

www.iana.org/assignments/bgp-well-known-communities

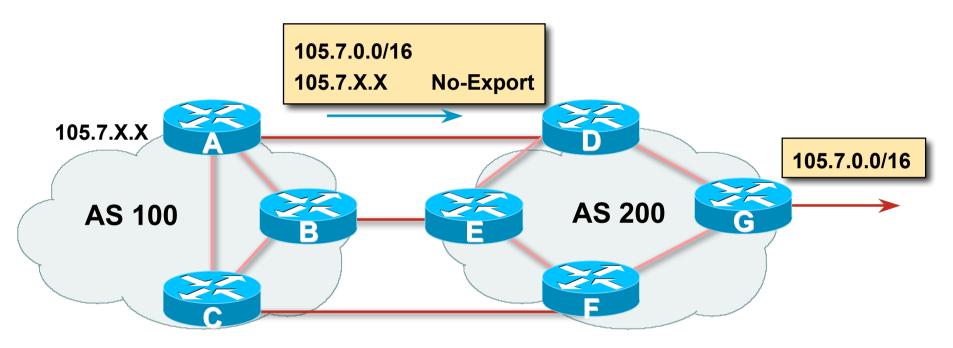
no-export
 65535:65281
 do not advertise to any eBGP peers

no-advertise
 do not advertise to any BGP peer

no-export-subconfed 65535:65283
 do not advertise outside local AS (only used with confederations)

no-peer
 do not advertise to bi-lateral peers (RFC3765)

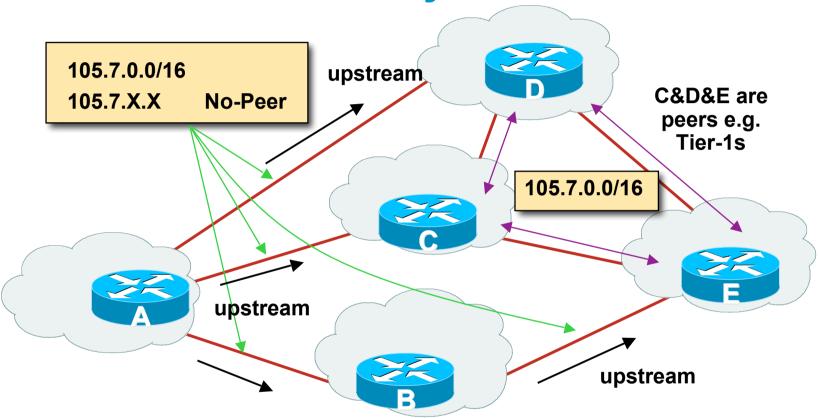
No-Export Community



- AS100 announces aggregate and subprefixes
 Intention is to improve loadsharing by leaking subprefixes
- Subprefixes marked with no-export community
- Router G in AS200 does not announce prefixes with no-export community set

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No-Peer Community



 Sub-prefixes marked with no-peer community are not sent to bi-lateral peers

They are only sent to upstream providers

Community Implementation details

Community is an optional attribute

Some implementations send communities to iBGP peers by default, some do not

Some implementations send communities to eBGP peers by default, some do not

Being careless can lead to community "confusion"

ISPs need consistent community policy within their own networks

And they need to inform peers, upstreams and customers about their community expectations

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BGP Path Selection Algorithm

Why Is This the Best Path?

BGP Path Selection Algorithm for IOS Part One

- Do not consider path if no route to next hop
- Do not consider iBGP path if not synchronised (Cisco IOS only)
- Highest weight (local to router)
- Highest local preference (global within AS)
- Prefer locally originated route
- Shortest AS path

BGP Path Selection Algorithm for IOS Part Two

- Lowest origin codeIGP < EGP < incomplete
- Lowest Multi-Exit Discriminator (MED)

If bgp deterministic-med, order the paths before comparing

(BGP spec does not specify in which order the paths should be compared. This means best path depends on order in which the paths are compared.)

If bgp always-compare-med, then compare for all paths otherwise MED only considered if paths are from the same AS (default)

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BGP Path Selection Algorithm for IOS Part Three

- Prefer eBGP path over iBGP path
- Path with lowest IGP metric to next-hop
- Lowest router-id (originator-id for reflected routes)
- Shortest Cluster-List
 Client must be aware of Route Reflector attributes!
- Lowest neighbour IP address

BGP Path Selection Algorithm

In multi-vendor environments:

Make sure the path selection processes are understood for each brand of equipment

Each vendor has slightly different implementations, extra steps, extra features, etc

Watch out for possible MED confusion

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Applying Policy with BGP

Controlling Traffic Flow & Traffic Engineering

Applying Policy in BGP: Why?

- Network operators rarely "plug in routers and go"
- External relationships:

Control who they peer with

Control who they give transit to

Control who they get transit from

Traffic flow control:

Efficiently use the scarce infrastructure resources (external link load balancing)

Congestion avoidance

Terminology: Traffic Engineering

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Applying Policy in BGP: How?

Policies are applied by:

Setting BGP attributes (local-pref, MED, AS-PATH, community), thereby influencing the path selection process

Advertising or Filtering prefixes

Advertising or Filtering prefixes according to ASN and AS-PATHs

Advertising or Filtering prefixes according to Community membership

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Applying Policy with BGP: Tools

• Most implementations have tools to apply policies to BGP:

Prefix manipulation/filtering

AS-PATH manipulation/filtering

Community Attribute setting and matching

 Implementations also have policy language which can do various match/set constructs on the attributes of chosen BGP routes

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Extending BGP

- Documented in RFC2842
- Capabilities parameters passed in BGP open message
- Unknown or unsupported capabilities will result in NOTIFICATION message
- Codes:

0 to 63 are assigned by IANA by IETF consensus 64 to 127 are assigned by IANA "first come first served" 128 to 255 are vendor specific

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Current capabilities are:

0	Reserved	[RFC3392]
1	Multiprotocol Extensions for BGP-4	[RFC4760]
2	Route Refresh Capability for BGP-4	[RFC2918]
3	Cooperative Route Filtering Capability	[ID]
4	Multiple routes to a destination capability	[RFC3107]
64	Graceful Restart Capability	[RFC4724]
65	Support for 4 octet ASNs	[RFC4893]
66	Deprecated 2003-03-06	
67	Support for Dynamic Capability	[ID]
68	Multisession BGP	[ID]

See www.iana.org/assignments/capability-codes

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Multiprotocol extensions

This is a whole different world, allowing BGP to support more than IPv4 unicast routes

Examples include: v4 multicast, IPv6, v6 multicast, VPNs Another tutorial (or many!)

- Route refresh is a well known scaling technique covered shortly
- 32-bit ASNs have recently arrived
- The other capabilities are still in development or not widely implemented or deployed yet

BGP for Internet Service Providers

- BGP Basics
- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



BGP Scaling Techniques

BGP Scaling Techniques

How does a service provider:

Scale the iBGP mesh beyond a few peers?

Implement new policy without causing flaps and route churning?

Keep the network stable, scalable, as well as simple?

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BGP Scaling Techniques

- Route Refresh
- Route Reflectors
- Confederations



Dynamic Reconfiguration

Route Refresh

Route Refresh

- BGP peer reset required after every policy change Because the router does not store prefixes which are rejected by policy
- Hard BGP peer reset:

Terminates BGP peering & Consumes CPU Severely disrupts connectivity for all networks

Soft BGP peer reset (or Route Refresh):

BGP peering remains active

Impacts only those prefixes affected by policy change

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Route Refresh Capability

- Facilitates non-disruptive policy changes
- For most implementations, no configuration is needed
 Automatically negotiated at peer establishment
- No additional memory is used
- Requires peering routers to support "route refresh capability" – RFC2918

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Dynamic Reconfiguration

- Use Route Refresh capability if supported find out from the BGP neighbour status display Non-disruptive, "Good For the Internet"
- If not supported, see if implementation has a workaround
- Only hard-reset a BGP peering as a last resort

Consider the impact to be equivalent to a router reboot



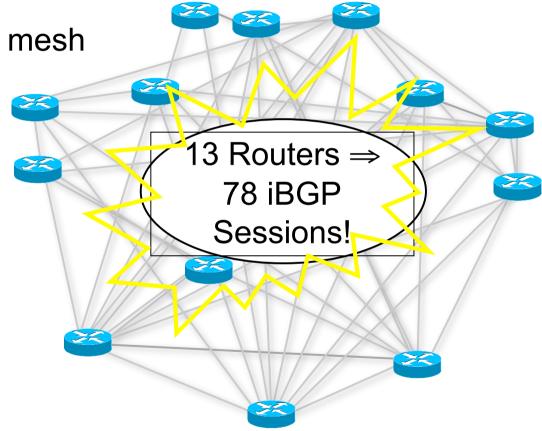
Route Reflectors

Scaling the iBGP mesh

Scaling iBGP mesh

Avoid ½n(n-1) iBGP mesh

n=1000 ⇒ nearly half a million ibgp sessions!

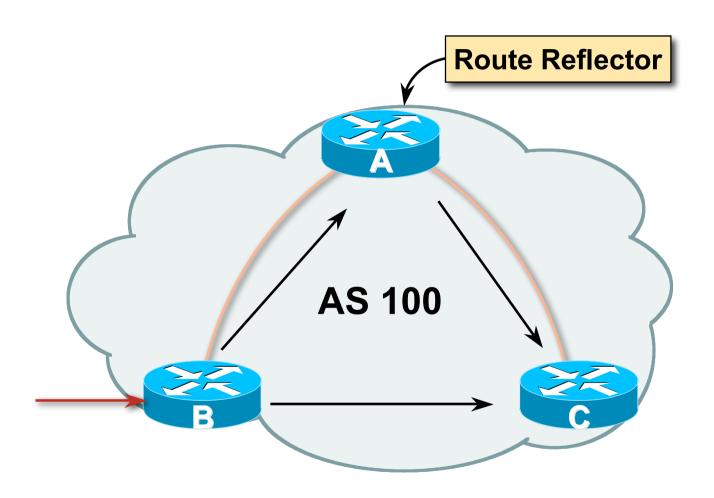


Two solutions

Route reflector – simpler to deploy and run

Confederation – more complex, has corner case advantages

Route Reflector: Principle



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Route Reflector

 Reflector receives path from clients and non-clients

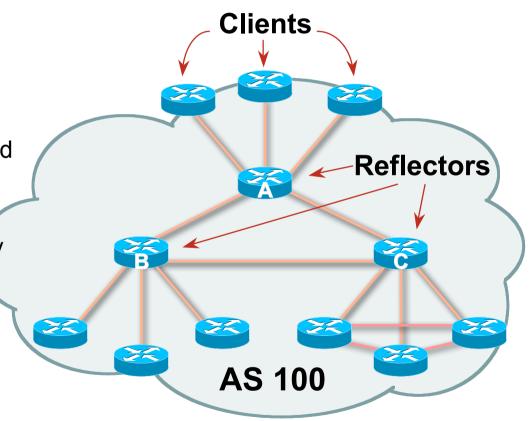
Selects best path

 If best path is from client, reflect to other clients and non-clients

 If best path is from non-client, reflect to clients only

Non-meshed clients

Described in RFC4456



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Route Reflector: Topology

- Divide the backbone into multiple clusters
- At least one route reflector and few clients per cluster
- Route reflectors are fully meshed
- Clients in a cluster could be fully meshed
- Single IGP to carry next hop and local routes

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Route Reflector: Loop Avoidance

Originator_ID attribute

Carries the RID of the originator of the route in the local AS (created by the RR)

Cluster list attribute

The local cluster-id is added when the update is sent by the RR Best to set cluster-id is from router-id (address of loopback) (Some ISPs use their own cluster-id assignment strategy – but needs to be well documented!)

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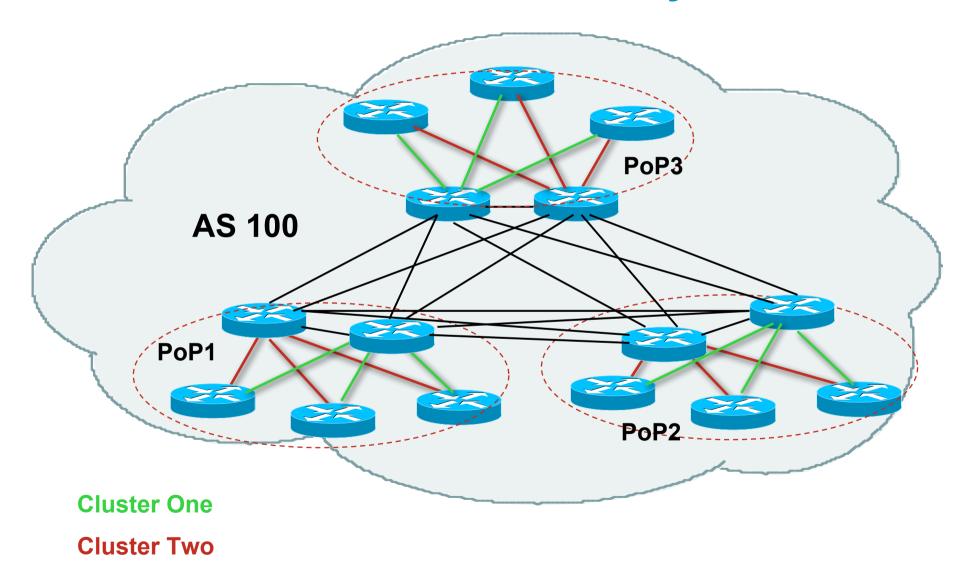
Route Reflector: Redundancy

• Multiple RRs can be configured in the same cluster – not advised!

All RRs in the cluster must have the same cluster-id (otherwise it is a different cluster)

- A router may be a client of RRs in different clusters
 - Common today in ISP networks to overlay two clusters redundancy achieved that way
 - → Each client has two RRs = redundancy

Route Reflector: Redundancy



Route Reflector: Benefits

- Solves iBGP mesh problem
- Packet forwarding is not affected
- Normal BGP speakers co-exist
- Multiple reflectors for redundancy
- Easy migration
- Multiple levels of route reflectors

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Route Reflector: Deployment

• Where to place the route reflectors?

Always follow the physical topology!

This will guarantee that the packet forwarding won't be affected

Typical ISP network:

PoP has two core routers

Core routers are RR for the PoP

Two overlaid clusters

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Route Reflector: Migration

Typical ISP network:

Core routers have fully meshed iBGP
Create further hierarchy if core mesh too big
Split backbone into regions

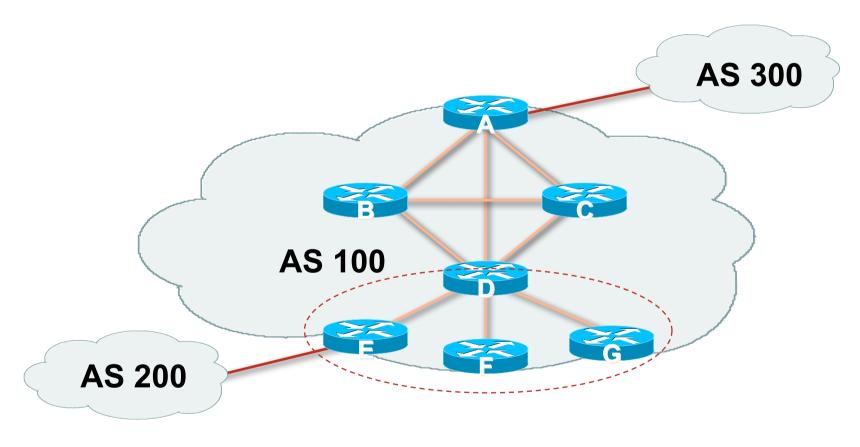
Configure one cluster pair at a time

Eliminate redundant iBGP sessions

Place maximum one RR per cluster

Easy migration, multiple levels

Route Reflector: Migration



 Migrate small parts of the network, one part at a time

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BGP Confederations

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Confederations

Divide the AS into sub-AS

eBGP between sub-AS, but some iBGP information is kept

Preserve NEXT_HOP across the sub-AS (IGP carries this information)

Preserve LOCAL_PREF and MED

- Usually a single IGP
- Described in RFC5065

Confederations (Cont.)

Visible to outside world as single AS – "Confederation Identifier"

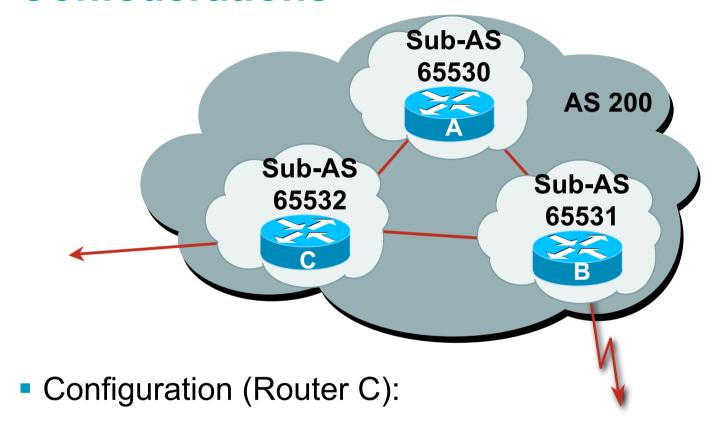
Each sub-AS uses a number from the private AS range (64512-65534)

iBGP speakers in each sub-AS are fully meshed

The total number of neighbours is reduced by limiting the full mesh requirement to only the peers in the sub-AS

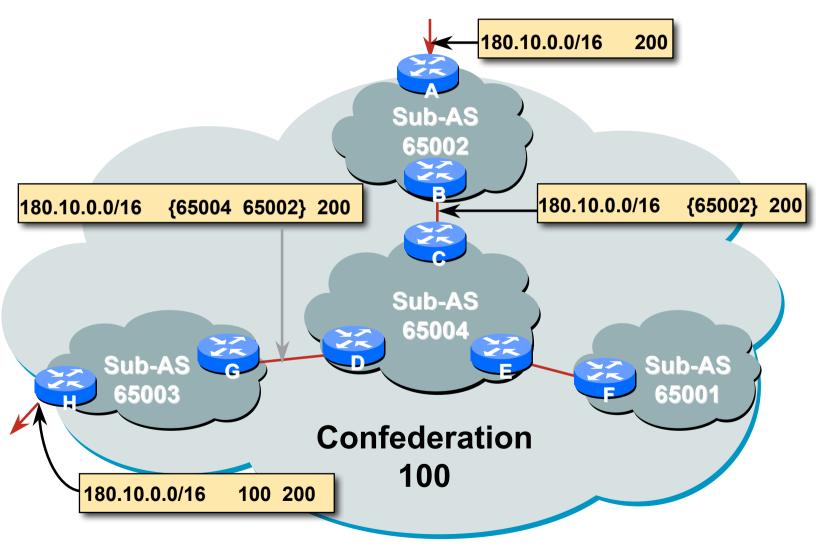
Can also use Route-Reflector within sub-AS

Confederations



bgp confederation identifier 200 bgp confederation peers 65530 65531 neighbor 141.153.12.1 remote-as 65530 neighbor 141.153.17.2 remote-as 65531

Confederations: AS-Sequence



Route Propagation Decisions

Same as with "normal" BGP:

From peer in same sub-AS → only to external peers

From external peers → to all neighbors

"External peers" refers to

Peers outside the confederation

Peers in a different sub-AS

Preserve LOCAL_PREF, MED and NEXT_HOP

RRs or Confederations

	Internet Connectivity	Multi-Level Hierarchy	Policy Control	Scalability	Migration Complexity
Confederations	Anywhere in the Network	Yes	Yes	Medium	Medium to High
Route Reflectors	Anywhere in the Network	Yes	Yes	Very High	Very Low

Most new service provider networks now deploy Route Reflectors from Day One

More points about Confederations

 Can ease "absorbing" other ISPs into you ISP – e.g., if one ISP buys another

Or can use AS masquerading feature available in some implementations to do a similar thing

 Can use route-reflectors with confederation sub-AS to reduce the sub-AS iBGP mesh



Route Flap Damping

Network Stability for the 1990s

Network Instability for the 21st Century!

Route Flap Damping

- For many years, Route Flap Damping was a strongly recommended practice
- Now it is strongly discouraged as it appears to cause far greater network instability than it cures
- But first, the theory...

Route Flap Damping

Route flap

Going up and down of path or change in attribute

BGP WITHDRAW followed by UPDATE = 1 flap

eBGP neighbour going down/up is NOT a flap

Ripples through the entire Internet

Wastes CPU

Damping aims to reduce scope of route flap propagation

Route Flap Damping (continued)

Requirements

Fast convergence for normal route changes

History predicts future behaviour

Suppress oscillating routes

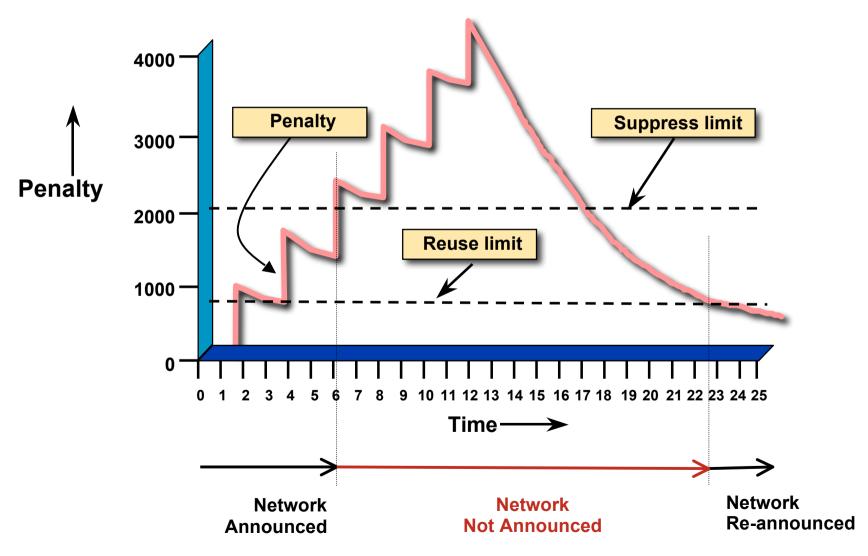
Advertise stable routes

Implementation described in RFC 2439

Operation

- Add penalty (1000) for each flap
 Change in attribute gets penalty of 500
- Exponentially decay penalty half life determines decay rate
- Penalty above suppress-limit do not advertise route to BGP peers
- Penalty decayed below reuse-limit
 re-advertise route to BGP peers
 penalty reset to zero when it is half of reuse-limit

Operation



Operation

- Only applied to inbound announcements from eBGP peers
- Alternate paths still usable
- Controllable by at least:

Half-life

reuse-limit

suppress-limit

maximum suppress time

Configuration

- Implementations allow various policy control with flap damping
 - Fixed damping, same rate applied to all prefixes
 - Variable damping, different rates applied to different ranges of prefixes and prefix lengths

Route Flap Damping History

- First implementations on the Internet by 1995
- Vendor defaults too severe

RIPE Routing Working Group recommendations in ripe-178, ripe-210, and ripe-229

http://www.ripe.net/ripe/docs

But many ISPs simply switched on the vendors' default values without thinking

Serious Problems:

 "Route Flap Damping Exacerbates Internet Routing Convergence"

Zhuoqing Morley Mao, Ramesh Govindan, George Varghese & Randy H. Katz, August 2002

- "What is the sound of one route flapping?"
 Tim Griffin, June 2002
- Various work on routing convergence by Craig Labovitz and Abha Ahuja a few years ago
- "Happy Packets"
 Closely related work by Randy Bush et al.

Problem 1:

One path flaps:

BGP speakers pick next best path, announce to all peers, flap counter incremented

Those peers see change in best path, flap counter incremented

After a few hops, peers see multiple changes simply caused by a single flap → prefix is suppressed

Problem 2:

 Different BGP implementations have different transit time for prefixes

Some hold onto prefix for some time before advertising Others advertise immediately

 Race to the finish line causes appearance of flapping, caused by a simple announcement or path change → prefix is suppressed

Solution:

- Do NOT use Route Flap Damping whatever you do!
- RFD will unnecessarily impair access to your network and to the Internet
- More information contained in RIPE Routing Working Group recommendations:

www.ripe.net/ripe/docs/ripe-378.[pdf,html,txt]

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Service Provider use of Communities

Some examples of how ISPs make life easier for themselves

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BGP Communities

- Another ISP "scaling technique"
- Prefixes are grouped into different "classes" or communities within the ISP network
- Each community means a different thing, has a different result in the ISP network

BGP Communities

 Communities are generally set at the edge of the ISP network

Customer edge: customer prefixes belong to different communities depending on the services they have purchased

Internet edge: transit provider prefixes belong to difference communities, depending on the loadsharing or traffic engineering requirements of the local ISP, or what the demands from its BGP customers might be

Two simple examples follow to explain the concept

- This demonstrates how communities might be used at the customer edge of an ISP network
- ISP has three connections to the Internet:
 - IXP connection, for local peers
 - Private peering with a competing ISP in the region
 - Transit provider, who provides visibility to the entire Internet
- Customers have the option of purchasing combinations of the above connections

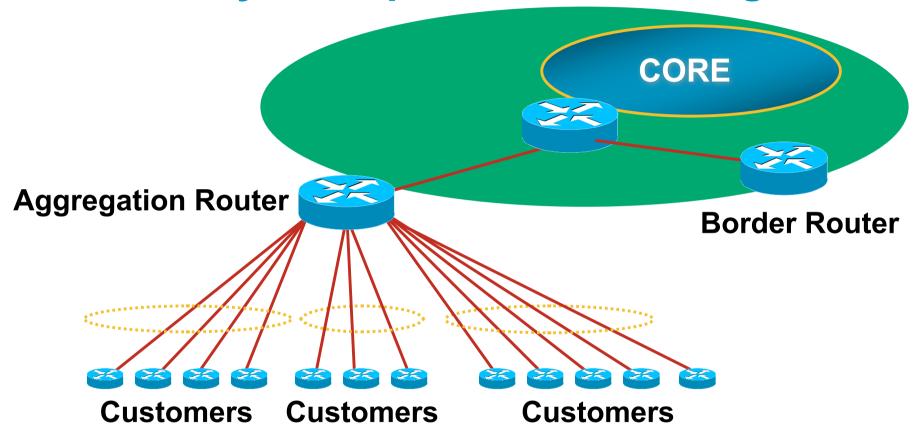
Community assignments:

IXP connection: community 100:2100

Private peer: community 100:2200

 Customer who buys local connectivity (via IXP) is put in community 100:2100

- Customer who buys peer connectivity is put in community 100:2200
- Customer who wants both IXP and peer connectivity is put in 100:2100 and 100:2200
- Customer who wants "the Internet" has no community set
 We are going to announce his prefix everywhere



Communities set at the aggregation router where the prefix is injected into the ISP's iBGP

- No need to alter filters at the network border when adding a new customer
- New customer simply is added to the appropriate community

Border filters already in place take care of announcements

⇒ Ease of operation!

Community Example: Internet Edge

- This demonstrates how communities might be used at the peering edge of an ISP network
- ISP has four types of BGP peers:

Customer

IXP peer

Private peer

Transit provider

- The prefixes received from each can be classified using communities
- Customers can opt to receive any or all of the above

Community Example: Internet Edge

Community assignments:

Customer prefix: community 100:3000

IXP prefix: community 100:3100

Private peer prefix: community 100:3200

- BGP customer who buys local connectivity gets 100:3000
- BGP customer who buys local and IXP connectivity receives community 100:3000 and 100:3100
- BGP customer who buys full peer connectivity receives community 100:3000, 100:3100, and 100:3200
- Customer who wants "the Internet" gets everything Gets default route originated by aggregation router
 Or pays money to get all 220k prefixes

Community Example: Internet Edge

 No need to create customised filters when adding customers

Border router already sets communities

Installation engineers pick the appropriate community set when establishing the customer BGP session

⇒ Ease of operation!

Community Example – Summary

- Two examples of customer edge and internet edge can be combined to form a simple community solution for ISP prefix policy control
- More experienced operators tend to have more sophisticated options available

Advice is to start with the easy examples given, and then proceed onwards as experience is gained

ISP BGP Communities

There are no recommended ISP BGP communities apart from

RFC1998

The five standard communities

www.iana.org/assignments/bgp-well-known-communities

Efforts have been made to document from time to time

totem.info.ucl.ac.be/publications/papers-elec-versions/draft-quoitin-bgp-comm-survey-00.pdf

But so far... nothing more... ⊗

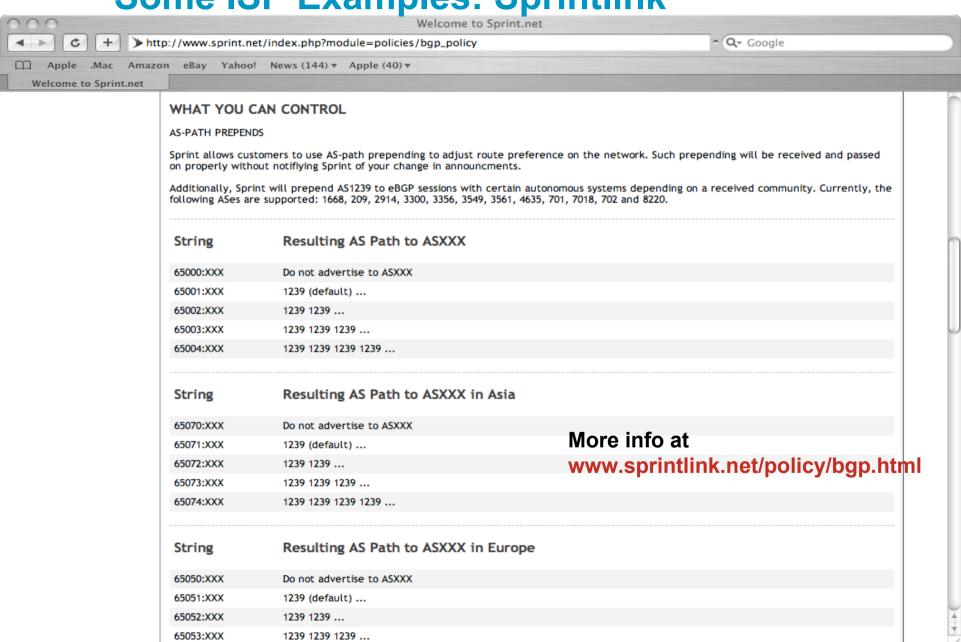
Collection of ISP communities at www.onesc.net/communities

ISP policy is usually published

On the ISP's website

Referenced in the AS Object in the IRR

Some ISP Examples: Sprintlink



Some ISP Examples AAPT

- Australian ISP
- Run their own Routing Registry
 Whois.connect.com.au
- Offer 6 different communities to customers to aid with their traffic engineering

Some ISP Examples AAPT

```
AS2764
aut-num:
              ASN-CONNECT-NET
as-name:
descr:
              AAPT Limited
admin-c:
              CNO2-AP
tech-c:
              CNO2-AP
              Community support definitions
remarks:
remarks:
remarks:
              Community Definition
remarks:
              2764:2 Don't announce outside local POP
remarks:
remarks:
              2764:4 Lower local preference by 15
remarks:
              2764:5 Lower local preference by 5
remarks:
              2764:6 Announce to customers and all peers
                           (incl int'l peers), but not transit
remarks:
              2764:7 Announce to customers only
remarks:
              2764:14 Announce to AANX
notify:
              routing@connect.com.au
mnt-by:
              CONNECT-AU
              nobody@connect.com.au 20050225
changed:
              CCAIR
source:
```

More at http://info.connect.com.au/docs/routing/general/multi-faq.shtml#q13

Some ISP Examples Verizon Business EMEA

- Verizon Business' European operation
- Permits customers to send communities which determine

local preferences within Verizon Business' network

Reachability of the prefix

How the prefix is announced outside of Verizon Business' network

Some ISP Examples Verizon Business Europe

```
aut-num: AS702
descr: Verizon Business EMEA - Commercial IP service provider in Eur
remarks: VzBi uses the following communities with its customers:
        702:80
                  Set Local Pref 80 within AS702
        702:120
                  Set Local Pref 120 within AS702
        702:20
                  Announce only to VzBi AS'es and VzBi customers
        702:30
                  Keep within Europe, don't announce to other VzBi AS
        702:1
                  Prepend AS702 once at edges of VzBi to Peers
        702:2
                  Prepend AS702 twice at edges of VzBi to Peers
        702:3
                  Prepend AS702 thrice at edges of VzBi to Peers
        Advanced communities for customers
        702:7020
                  Do not announce to AS702 peers with a scope of
                  National but advertise to Global Peers, European
                  Peers and VzBi customers.
        702:7001 Prepend AS702 once at edges of VzBi to AS702
                  peers with a scope of National.
        702:7002
                  Prepend AS702 twice at edges of VzBi to AS702
                  peers with a scope of National.
(more)
```

Some ISP Examples VzBi Europe

```
(more)
         702:7003 Prepend AS702 thrice at edges of VzBi to AS702
                  peers with a scope of National.
         702:8020 Do not announce to AS702 peers with a scope of
                  European but advertise to Global Peers, National
                  Peers and VzBi customers.
         702:8001 Prepend AS702 once at edges of VzBi to AS702
                 peers with a scope of European.
         702:8002 Prepend AS702 twice at edges of VzBi to AS702
                  peers with a scope of European.
         702:8003 Prepend AS702 thrice at edges of VzBi to AS702
                 peers with a scope of European.
         Additional details of the VzBi communities are located at:
         http://www.verizonbusiness.com/uk/customer/bqp/
mnt-by: WCOM-EMEA-RICE-MNT
source: RIPE
```

Some ISP Examples BT Ignite

One of the most comprehensive community lists around

Seems to be based on definitions originally used in Tiscali's network

whois -h whois.ripe.net AS5400 reveals all

 Extensive community definitions allow sophisticated traffic engineering by customers

Some ISP Examples BT Ignite

aut-num:	AS5400	
descr:	BT Ignite European Backbone	
remarks:		
remarks:	Community to	Community to
remarks:	Not announce To peer:	AS prepend 5400
remarks:		
remarks:	5400:1000 All peers & Transits	5400:2000
remarks:		
remarks:	5400:1500 All Transits	5400:2500
remarks:	5400:1501 Sprint Transit (AS1239)	5400:2501
remarks:	5400:1502 SAVVIS Transit (AS3561)	5400:2502
remarks:	5400:1503 Level 3 Transit (AS3356)	5400:2503
remarks:	5400:1504 AT&T Transit (AS7018)	5400:2504
remarks:	5400:1506 GlobalCrossing Trans (AS3549)	5400:2506
remarks:		
remarks:	5400:1001 Nexica (AS24592)	5400:2001
remarks:	5400:1002 Fujitsu (AS3324)	5400:2002
remarks:	5400:1004 C&W EU (1273)	5400:2004
<snip></snip>		
notify:	notify@eu.bt.net And mar	NV.
mnt-by:	CTP-MNT Allu IIIal	
source:	RIPE many mo	rei

Some ISP Examples Level 3

- Highly detailed AS object held on the RIPE Routing Registry
- Also a very comprehensive list of community definitions

whois -h whois.ripe.net AS3356 reveals all

Some ISP Examples Level 3

```
AS3356
aut-num:
             Level 3 Communications
descr:
<snip>
remarks:
             customer traffic engineering communities - Suppression
remarks:
remarks:
             64960:XXX - announce to AS XXX if 65000:0
remarks:
remarks:
             65000:0 - announce to customers but not to peers
remarks:
             65000:XXX - do not announce at peerings to AS XXX
remarks:
remarks:
             customer traffic engineering communities - Prepending
remarks:
             65001:0 - prepend once to all peers
remarks:
remarks:
             65001:XXX - prepend once at peerings to AS XXX
<snip>
             3356:70 - set local preference to 70
remarks:
remarks:
          3356:80 - set local preference to 80
remarks: 3356:90 - set local preference to 90
remarks:
             3356:9999 - blackhole (discard) traffic
<snip>
mnt-by:
             LEVEL3-MNT
                                                  And many
source:
             RIPE
                                                 many more!
```

BGP for Internet Service Providers

- BGP Basics
- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



Deploying BGP in an ISP Network

Okay, so we've learned all about BGP now; how do we use it on our network??

Deploying BGP

- The role of IGPs and iBGP
- Aggregation
- Receiving Prefixes
- Configuration Tips



The role of IGP and iBGP

Ships in the night?

Or

Good foundations?

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BGP versus OSPF/ISIS

Internal Routing Protocols (IGPs)

examples are ISIS and OSPF

used for carrying infrastructure addresses

NOT used for carrying Internet prefixes or customer prefixes

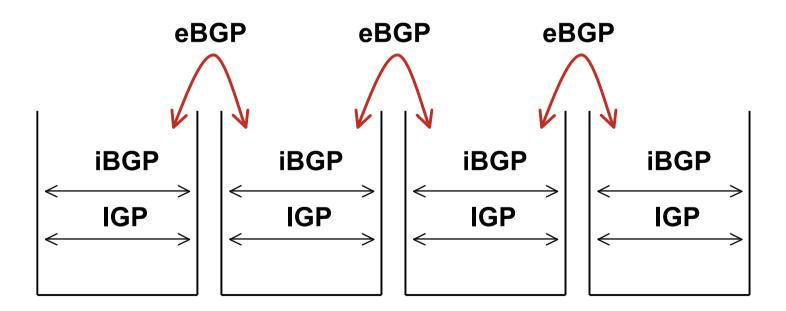
design goal is to minimise number of prefixes in IGP to aid scalability and rapid convergence

BGP versus OSPF/ISIS

- BGP used internally (iBGP) and externally (eBGP)
- iBGP used to carry some/all Internet prefixes across backbone customer prefixes
- eBGP used to exchange prefixes with other ASes implement routing policy

BGP/IGP model used in ISP networks

Model representation



BGP versus OSPF/ISIS

DO NOT:

distribute BGP prefixes into an IGP distribute IGP routes into BGP use an IGP to carry customer prefixes

YOUR NETWORK WILL NOT SCALE

Injecting prefixes into iBGP

- Use iBGP to carry customer prefixes
 Don't ever use IGP
- Point static route to customer interface
- Enter network into BGP process
 - Ensure that implementation options are used so that the prefix always remains in iBGP, regardless of state of interface
 - i.e. avoid iBGP flaps caused by interface flaps



Aggregation

Quality or Quantity?

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Aggregation

- Aggregation means announcing the address block received from the RIR to the other ASes connected to your network
- Subprefixes of this aggregate may be:
 - Used internally in the ISP network
 - Announced to other ASes to aid with multihoming
- Unfortunately too many people are still thinking about class Cs, resulting in a proliferation of /24s in the Internet routing table

Aggregation

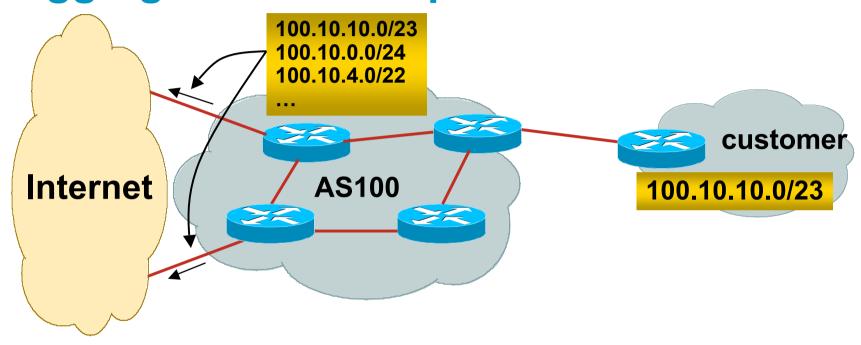
- Address block should be announced to the Internet as an aggregate
- Subprefixes of address block should NOT be announced to Internet unless special circumstances (more later)
- Aggregate should be generated internally Not on the network borders!

Announcing an Aggregate

- ISPs who don't and won't aggregate are held in poor regard by community
- Registries publish their minimum allocation size Anything from a /20 to a /22 depending on RIR Different sizes for different address blocks
- No real reason to see anything longer than a /22 prefix in the Internet

BUT there are currently >124000 /24s!

Aggregation – Example



- Customer has /23 network assigned from AS100's /19 address block
- AS100 announces customers' individual networks to the Internet

Aggregation – Bad Example

Customer link goes down

Their /23 network becomes unreachable

/23 is withdrawn from AS100's iBGP

 Their ISP doesn't aggregate its /19 network block

/23 network withdrawal announced to peers

starts rippling through the Internet

added load on all Internet backbone routers as network is removed from routing table

Customer link returns

Their /23 network is now visible to their ISP

Their /23 network is readvertised to peers

Starts rippling through Internet

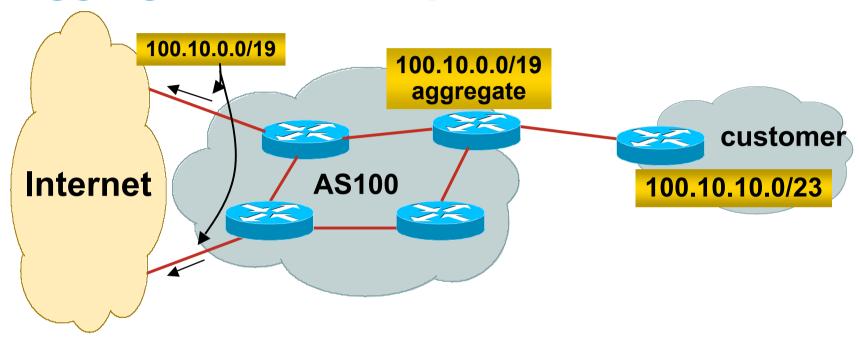
Load on Internet backbone routers as network is reinserted into routing table

Some ISP's suppress the flaps

Internet may take 10-20 min or longer to be visible

Where is the Quality of Service???

Aggregation – Example



- Customer has /23 network assigned from AS100's /19 address block
- AS100 announced /19 aggregate to the Internet

Aggregation – Good Example

- Customer link goes down their /23 network becomes unreachable
 /23 is withdrawn from AS100's iBGP
- /19 aggregate is still being announced
 no BGP hold down problems
 no BGP propagation delays
 no damping by other ISPs

- Customer link returns
- Their /23 network is visible again

The /23 is re-injected into AS100's iBGP

- The whole Internet becomes visible immediately
- Customer has Quality of Service perception

Aggregation – Summary

Good example is what everyone should do!

Adds to Internet stability

Reduces size of routing table

Reduces routing churn

Improves Internet QoS for everyone

Bad example is what too many still do!

Why? Lack of knowledge?

Laziness?

The Internet Today (November 2007)

Current Internet Routing Table Statistics

BGP Routing Table Entries	237854
Prefixes after maximum aggregation	122667
Unique prefixes in Internet	116358
Prefixes smaller than registry alloc	110699
/24s announced	124638
only 5727 /24s are from 192.0.0.0/8	
ASes in use	26654

"The New Swamp"

Swamp space is name used for areas of poor aggregation

The original swamp was 192.0.0.0/8 from the former class C block

Name given just after the deployment of CIDR

The new swamp is creeping across all parts of the Internet Not just RIR space, but "legacy" space too

"The New Swamp" RIR Space – February 1999

RIR blocks contribute 49393 prefixes or 88% of the Internet Routing Table

Block	Networks	Block	Networks	Block	Networks	Block	Networks
24/8	165	77/8	0	118/8	0	203/8	3622
41/8	0	78/8	0	119/8	0	204/8	3792
58/8	0	79/8	0	120/8	0	205/8	2584
59/8	0	80/8	0	121/8	0	206/8	3127
60/8	0	81/8	0	122/8	0	207/8	2723
61/8	3	82/8	0	123/8	0	208/8	2817
62/8	87	83/8	0	124/8	0	209/8	2574
63/8	20	84/8	0	125/8	0	210/8	617
64/8	0	<i>85/8</i>	0	126/8	0	211/8	0
65/8	0	86/8	0	189/8	0	212/8	717
66/8	0	87/8	0	190/8	0	213/8	1
67/8	0	88/8	0	192/8	6275	216/8	943
68/8	0	89/8	0	193/8	2390	217/8	0
69/8	0	90/8	0	194/8	2932	218/8	0
70/8	0	91/8	0	195/8	1338	219/8	0
71/8	0	96/8	0	196/8	513	220/8	0
72/8	0	97/8	0	198/8	4034	221/8	0
73/8	0	98/8	0	199/8	3495	222/8	0
74/8	0	99/8	0	200/8	1348		
75/8	0	116/8	0	201/8	0		
76/8	0	117/8	0	202/8	2276		

13/

"The New Swamp" RIR Space – February 2007

RIR blocks contribute 192490 prefixes or 90% of the Internet Routing Table

Block	Networks	Block	Networks	Block	Networks	Block	Networks
24/8	2930	77/8	1214	118/8	3	203/8	10459
41/8	288	78/8	8	119/8	3	204/8	5569
58/8	1097	79/8	2	120/8	3	205/8	2892
59/8	1152	80/8	2053	121/8	426	206/8	3857
60/8	604	81/8	1695	122/8	698	207/8	4331
61/8	2589	82/8	1564	123/8	534	208/8	4258
62/8	2193	83/8	1172	124/8	1340	209/8	5540
63/8	2967	84/8	1269	125/8	1554	210/8	4759
64/8	5501	85/8	1891	126/8	41	211/8	2733
65/8	3917	86/8	800	189/8	169	212/8	2900
66/8	6575	87/8	1157	190/8	1077	213/8	3052
67/8	2015	88/8	847	192/8	6927	216/8	6930
68/8	2770	89/8	1970	193/8	5704	217/8	2615
69/8	3702	90/8	105	194/8	4652	218/8	1561
70/8	1693	91/8	577	195/8	4279	219/8	1197
71/8	1188	96/8	8	196/8	1600	220/8	1988
72/8	2878	97/8	1	198/8	4748	221/8	894
73/8	273	98/8	3	199/8	4184	222/8	1241
74/8	1483	99/8	0	200/8	7482		
75/8	483	116/8	3 3	201/8	2927		
76/8	194	117/8	3	202/8	10529		

"The New Swamp" Summary

RIR space shows creeping deaggregation

It seems that an RIR /8 block averages around 5000 prefixes once fully allocated

So their existing 81 /8s will eventually cause 405000 prefix announcements

Food for thought:

Remaining 48 unallocated /8s and the 81 RIR /8s combined will cause:

645000 prefixes with 5000 prefixes per /8 density

774000 prefixes with 6000 prefixes per /8 density

Plus 12% due to "non RIR space deaggregation"

→ Routing Table size of 866880 prefixes

"The New Swamp" Summary

- Rest of address space is showing similar deaggregation too
- What are the reasons?Main justification is traffic engineering
- Real reasons are:

Lack of knowledge

Laziness

Deliberate & knowing actions

BGP Report (bgp.potaroo.net)

- 199336 total announcements in October 2006
- 129795 prefixes

After aggregating including full AS PATH info i.e. including each ASN's traffic engineering 35% saving possible

109034 prefixes

After aggregating by Origin AS i.e. ignoring each ASN's traffic engineering

10% saving possible

Deaggregation: The Excuses

- Traffic engineering causes 10% of the Internet Routing table
- Deliberate deaggregation causes 35% of the Internet Routing table

Efforts to improve aggregation

The CIDR Report

Initiated and operated for many years by Tony Bates
Now combined with Geoff Huston's routing analysis

www.cidr-report.org

Results e-mailed on a weekly basis to most operations lists around the world

Lists the top 30 service providers who could do better at aggregating

RIPE Routing WG aggregation recommendation

RIPE-399 — http://www.ripe.net/ripe/docs/ripe-399.html

Efforts to Improve Aggregation The CIDR Report

- Also computes the size of the routing table assuming ISPs performed optimal aggregation
- Website allows searches and computations of aggregation to be made on a per AS basis

Flexible and powerful tool to aid ISPs

Intended to show how greater efficiency in terms of BGP table size can be obtained without loss of routing and policy information

Shows what forms of origin AS aggregation could be performed and the potential benefit of such actions to the total table size

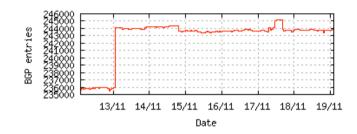
Very effectively challenges the traffic engineering excuse



Status Summary

Table History

Date	Prefixes	CIDR Aggregated
12-11-07	235665	148466
13-11-07	235821	154894
14-11-07	244197	155135
15-11-07	243680	155759
16-11-07	243479	157630
17-11-07	243672	157800
18-11-07	243699	158137
19-11-07	243780	158206



Plot: BGP Table Size

AS Summary

26760 Number of ASes in routing system 11268 Number of ASes announcing only one prefix

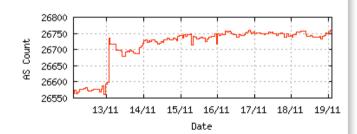
1966 Largest number of prefixes announced by an AS

AS4538: ERX-CERNET-BKB China Education and Research

Network Center

89036800 Largest address span announced by an AS (/32s)

AS721: DISA-ASNBLK - DoD Network Information Center



Plot: AS count

Plot: Average announcements per origin AS

Report: ASes ordered by originating address span Report: ASes ordered by transit address span

Report: Autonomous System number-to-name mapping (from Registry WHOIS data)



Aggregation Summary

The algorithm used in this report proposes aggregation only when there is a precise match using AS path so as to preserve traffic transit policies. Aggregation is also proposed across non-advertised address space ('holes').

--- 19Nov07 ---

ASnum	NetsNow	NetsAggr	NetGain	9/ Gain	^o Description	
Table	243774	158123	85651	35.1%	All ASes	
AS4538	1966	718	1248	63.5%	ERX-CERNET-BKB China Education and Research Network Center	
AS4755	1467	315	1152	78.5%	VSNL-AS Videsh Sanchar Nigam Ltd. Autonomous System	
AS9498	1050	69	981	93.4%	BBIL-AP BHARTI BT INTERNET LTD.	
AS4323	1373	473	900	65.5%	TWTC - Time Warner Telecom, Inc.	
AS11492	1173	422	751	64.0%	CABLEONE - CABLE ONE	
AS22773	819	74	745	91.0%	CCINET-2 - Cox Communications Inc.	
AS6478	1124	393	731	65.0%	ATT-INTERNET3 - AT&T WorldNet Services	
AS8151	1161	439	722	62.2%	Uninet S.A. de C.V.	
AS18566	1032	360	672	65.1%	COVAD - Covad Communications Co.	
AS19262	811	183	628	77.4%	VZGNI-TRANSIT - Verizon Internet Services Inc.	
AS17488	879	265	614	69.9%	HATHWAY-NET-AP Hathway IP Over Cable Internet	
AS15270	603	41	562	93.2%	AS-PAETEC-NET - PaeTec Communications, Inc.	
AS18101	611	74	537	87.9%	RIL-IDC Reliance Infocom Ltd Internet Data Centre,	
AS4134	1103	596	507	46.0%	CHINANET-BACKBONE No.31, Jin-rong Street	
AS7545	725	230	495		TPG-INTERNET-AP TPG Internet Pty Ltd	
AS7018	1518	1028	490		ATT-INTERNET4 - AT&T WorldNet Services	
AS6197	1031	554	477		BATI-ATL - BellSouth Network Solutions, Inc	
AS2386	1276	800	476		INS-AS - AT&T Data Communications Services	
AS4766	817	374	443		KIXS-AS-KR Korea Telecom	
AS4812	525	88	437		CHINANET-SH-AP China Telecom (Group)	
AS7011	982	574	408		FRONTIER-AND-CITIZENS - Frontier Communications of America, Inc.	
AS17676	503	126	377	75.0%	GIGAINFRA BB TECHNOLOGY Corp.	
AS4808	503	128	375	74.6%	CHINA169-BJ CNCGROUP IP network China169 Beijing Province Network	
AS9443	450	76	374	83.1%	INTERNETPRIMUS-AS-AP Primus Telecommunications	
AS19916	569	206	363	63.8%	ASTRUM-0001 - OLM LLC	



Top 20 Added Routes this week per Originating AS

```
Prefixes ASnum AS Description
     722 AS23577 ATM-MPLS-AS-KR Korea Telecom
     265 AS3464 ASC-NET - Alabama Supercomputer Network
     250 AS31283 FASTHOST-AS FastHost AS - Norwegian based ISP
     209 AS1659 ERX-TANET-ASN1 Tiawan Academic Network (TANet) Information Center
     203 AS11426 SCRR-11426 - Road Runner HoldCo LLC
     198 AS4668
                 LGNET-AS-KR LG CNS
     193 AS10154 USE-AS-KR ULSAN Office of Education
     178 AS2686 AT&T Global Network Services - EMEA
     176 AS16473 TNII - Bell South
     161 AS19548 ADELPHIA-AS2 - Road Runner HoldCo LLC
     149 AS4538 ERX-CERNET-BKB China Education and Research Network Center
     146 AS10316 ABACUS-NET-AS - Abacus America Inc.
     138 AS2819 GTSCZ GTS NOVERA (GTS CZ)
     137 AS1273
                 CW Cable & Wireless
     129 AS3561
                 SAVVIS - Savvis
     128 AS38394 GOESN-AS-KR Gyeonggido Seongnam Office of Education
     121 AS3300 BT-INFONET-EUROPE BT-Infonet-Europe
     120 AS237
                 MERIT-AS-14 - Merit Network Inc.
     115 AS23216 MEGADATOS S.A.
      93 AS376
                 RISO-AS - Reseau Interordinateurs Scientique Quebecois (RISO)
```

Top 20 Withdrawn Routes this week per Originating AS

Prefixes ASnum AS Description -91 AS7725 CCH-AS7 - Comcast Cable Communications Holdings, Inc -69 AS14359 ITS-USNET - Ideal Technology Solutions US Inc. -53 AS4274 ERX-AU-NET Assumption University -46 AS17917 ECLTELECOMM-AS-AP ECL TeleCommunication Ltd --No Registry Entry---40 AS7455 -39 AS8452 TEDATA TEDATA FR-RENATER Reseau National de telecommunications pour la Technologie -33 AS2200 -33 AS36728 EMERY-TELCOM-0000 - EMERY TELCOM GENESIS-AP Divixian.com Limited -31 AS9584 -29 AS3246 TDCSONG TDC Song AS11509 TIERZERO-AS11509 - Tierzero



More Specifics

A list of route advertisements that appear to be more specific than the original Class-based prefix mask, or more specific than the registry allocation size.

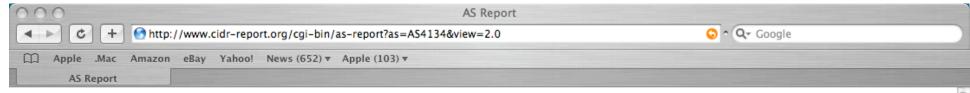
Top 20 ASes advertising more specific prefixes

More Specifics	Total Prefixes	ASnum	AS Description
1906	1966	AS4538	ERX-CERNET-BKB China Education and Research Network Center
1449	1467	AS4755	VSNL-AS Videsh Sanchar Nigam Ltd. Autonomous System
1247	1518	AS7018	ATT-INTERNET4 - AT&T WorldNet Services
1182	1373	AS4323	TWTC - Time Warner Telecom, Inc.
1180	1276	AS2386	INS-AS - AT&T Data Communications Services
1168	1173	AS11492	CABLEONE - CABLE ONE
1155	1161	AS8151	Uninet S.A. de C.V.
1124	1124	AS6478	ATT-INTERNET3 - AT&T WorldNet Services
1089	1089	AS9583	SIFY-AS-IN Sify Limited
1029	1050	AS9498	BBIL-AP BHARTI BT INTERNET LTD.
1023	1032	AS18566	COVAD - Covad Communications Co.
1010	1031	AS6197	BATI-ATL - BellSouth Network Solutions, Inc
973	982	AS7011	FRONTIER-AND-CITIZENS - Frontier Communications of America, Inc.
966	966	AS23577	ATM-MPLS-AS-KR Korea Telecom
879	879	AS17488	HATHWAY-NET-AP Hathway IP Over Cable Internet
831	846	AS20115	CHARTER-NET-HKY-NC - Charter Communications
811	1103	AS4134	CHINANET-BACKBONE No.31, Jin-rong Street
786	819	AS22773	CCINET-2 - Cox Communications Inc.
773	817	AS4766	KIXS-AS-KR Korea Telecom
771	811	AS19262	VZGNI-TRANSIT - Verizon Internet Services Inc.

Report: ASes ordered by number of more specific prefixes

Report: More Specific prefix list (by AS)

Report: More Specific prefix list (ordered by prefix)



Announced Prefixes

```
Rank AS Type Originate Addr Space (pfx) Transit Addr space (pfx) Description
4 AS4134 ORG+TRN Originate: 68834496 /5.96 Transit: 35830464 /6.91 CHINANET-BACKBONE No.31, Jin-rong Street
```

Aggregation Suggestions

This report does not take into account conditions local to each origin AS in terms of policy or traffic engineering requirements, so this is an approximate guideline as to aggregation possibilities.

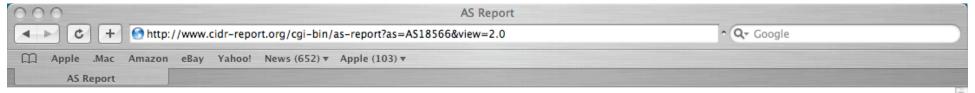
```
Rank AS
                   AS Name
                                                                Current Wthdw
                                                                                       Annce Redctn
                                                                                Aggte
  15 AS4134
                   CHINANET-BACKBONE No.31, Jin-rong Street
                                                                   1103
                                                                            665
                                                                                   158
                                                                                          596
                                                                                                 507 45.97%
Prefix
                      AS Path
                                                           Aggregation Suggestion
 58.30.0.0/15
                      12654 7018 4134
58.32.0.0/13
                      12654 7018 4134
58.40.0.0/15
                      12654 7018 4134
58.42.0.0/16
                      12654 3257 4134
                                           + Announce - aggregate of 58.42.0.0/17 (12654 3257 4134) and 58.42.128.0/17 (12654 3257
                                           - Withdrawn - aggregated with 58.42.128.0/17 (12654 3257 4134)
58.42.0.0/17
                      12654 3257 4134
                                           - Withdrawn - aggregated with 58.42.0.0/17 (12654 3257 4134)
58.42.128.0/17
                      12654 3257 4134
58.43.0.0/16
                      12654 7018 4134
58.44.0.0/14
                      12654 7018 4134
58.48.0.0/13
                      12654 7018 4134
58.48.0.0/13
                      12654 3257 4134
                                           + Announce - aggregate of 58.48.0.0/14 (12654 3257 4134) and 58.52.0.0/14 (12654 3257 4
58.48.0.0/14
                                           - Withdrawn - aggregated with 58.52.0.0/14 (12654 3257 4134)
                      12654 3257 4134
58.52.0.0/14
                      12654 3257 4134
                                           - Withdrawn - aggregated with 58.48.0.0/14 (12654 3257 4134)
58.56.0.0/15
                      12654 7018 4134
                                           + Announce - aggregate of 58.58.0.0/16 (12654 7018 4134) and 58.59.0.0/16 (12654 7018 4
58.58.0.0/15
                      12654 7018 4134
58.58.0.0/16
                                           - Withdrawn - aggregated with 58.59.0.0/16 (12654 7018 4134)
                      12654 7018 4134
58.59.0.0/17
                      12654 7018 4134
                                           - Withdrawn - aggregated with 58.59.128.0/17 (12654 7018 4134)
58.59.128.0/17
                      12654 7018 4134

    Withdrawn - aggregated with 58.59.0.0/17 (12654 7018 4134)

                                           - Withdrawn - matching aggregate 58.59.128.0/17 12654 7018 4134
58.59.128.0/19
                      12654 7018 4134
                                           - Withdrawn - matching aggregate 58.59.128.0/17 12654 7018 4134
58.59.160.0/19
                      12654 7018 4134
                                           - Withdrawn - matching aggregate 58.59.128.0/17 12654 7018 4134
58.59.192.0/19
                      12654 7018 4134
                                           - Withdrawn - matching aggregate 58.59.128.0/17 12654 7018 4134
58.59.224.0/19
                      12654 7018 4134
58.60.0.0/14
                      12654 7018 4134
                                           - Withdrawn - matching aggregate 58.60.0.0/14 12654 7018 4134
58.60.0.0/15
                      12654 7018 4134
58.62.0.0/15
                      12654 7018 4134
                                           - Withdrawn - matching aggregate 58.60.0.0/14 12654 7018 4134
58.66.0.0/15
                      12654 7018 4134
                                           + Announce - aggregate of 58.66.0.0/16 (12654 7018 4134) and 58.67.0.0/16 (12654 7018 4
58.66.0.0/17
                      12654 7018 4134

    Withdrawn - aggregated with 58.66.128.0/17 (12654 7018 4134)

58.66.128.0/18
                      12654 7018 4134
                                           - Withdrawn - aggregated with 58.66.192.0/18 (12654 7018 4134)
58.66.192.0/18
                      12654 7018 4134
                                           - Withdrawn - aggregated with 58.66.128.0/18 (12654 7018 4134)
                                           - Withdrawn - aggregated with 58.67.128.0/17 (12654 7018 4134)
58.67.0.0/17
                      12654 7018 4134
                                           - Withdrawn - aggregated with 58.67.0.0/17 (12654 7018 4134)
58.67.128.0/17
                      12654 7018 4134
 58.82.0.0/17
                      12654 7018 4134
```



Announced Prefixes

```
Rank AS Type Originate Addr Space (pfx) Transit Addr space (pfx) Description
147 AS18566 ORIGIN Originate: 2296320 /10.87 Transit: 0 /0.00 COVAD - Covad Communications Co.
```

Aggregation Suggestions

This report does not take into account conditions local to each origin AS in terms of policy or traffic engineering requirements, so this is an approximate guideline as to aggregation possibilities.

```
Rank AS
                   AS Name
                                                                         Wthdw
                                                                                Aggte
                                                                                      Annce Redctn
                   COVAD - Covad Communications Co.
  10 AS18566
                                                                   1032
                                                                           834
                                                                                  162
                                                                                         360
                                                                                                672 65.12%
Prefix
                      AS Path
                                                           Aggregation Suggestion
 64.105.0.0/16
                      12654 3257 2828 18566
 64.105.0.0/23
                      12654 3333 3356 18566
 64.105.4.0/22
                      12654 3333 3356 18566 + Announce - aggregate of 64.105.4.0/23 (12654 3333 3356 18566) and 64.105.6.0/23 (126
64.105.4.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.6.0/23 (12654 3333 3356 18566)
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.4.0/23 (12654 3333 3356 18566)
64.105.6.0/23
 64.105.8.0/22
                      12654 3333 3356 18566 + Announce - aggregate of 64.105.8.0/23 (12654 3333 3356 18566) and 64.105.10.0/23 (12
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.10.0/23 (12654 3333 3356 18566)
64.105.8.0/23
64.105.10.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.8.0/23 (12654 3333 3356 18566)
 64.105.14.0/23
                      12654 3333 3356 18566
 64.105.16.0/20
                      12654 3333 3356 18566 + Announce - aggregate of 64.105.16.0/21 (12654 3333 3356 18566) and 64.105.24.0/21 (1
64.105.16.0/24
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.17.0/24 (12654 3333 3356 18566)
64.105.17.0/24
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.16.0/24 (12654 3333 3356 18566)
64.105.18.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.16.0/23 (12654 3333 3356 18566)
64.105.20.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.22.0/23 (12654 3333 3356 18566)
 64.105.22.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.20.0/23 (12654 3333 3356 18566)
64.105.24.0/21
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.16.0/21 (12654 3333 3356 18566)
 64.105.32.0/20
                      12654 3333 3356 18566 + Announce - aggregate of 64.105.32.0/21 (12654 3333 3356 18566) and 64.105.40.0/21 (1
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.40.0/21 (12654 3333 3356 18566)
64.105.32.0/21
64.105.40.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.42.0/23 (12654 3333 3356 18566)
64.105.42.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.40.0/23 (12654 3333 3356 18566)
64.105.44.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.46.0/23 (12654 3333 3356 18566)
64.105.46.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.44.0/23 (12654 3333 3356 18566)
 64.105.48.0/21
                      12654 3333 3356 18566 + Announce - aggregate of 64.105.48.0/22 (12654 3333 3356 18566) and 64.105.52.0/22 (1
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.50.0/23 (12654 3333 3356 18566)
64.105.48.0/23
64.105.50.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.48.0/23 (12654 3333 3356 18566)
 64.105.52.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.54.0/23 (12654 3333 3356 18566)
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.52.0/23 (12654 3333 3356 18566)
64.105.54.0/23
 64.105.56.0/22
                      12654 3333 3356 18566 + Announce - aggregate of 64.105.56.0/23 (12654 3333 3356 18566) and 64.105.58.0/23 (1
64.105.56.0/23
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.58.0/23 (12654 3333 3356 18566)
                      12654 3333 3356 18566 - Withdrawn - aggregated with 64.105.56.0/23 (12654 3333 3356 18566)
 64.105.58.0/23
 64.105.60.0/23
                      12654 7018 3356 18566
```

Importance of Aggregation

Size of routing table

Memory is no longer a problem

Routers can be specified to carry 1 million prefixes

Convergence of the Routing System

This is a problem

Bigger table takes longer for CPU to process

BGP updates take longer to deal with

BGP Instability Report tracks routing system update activity

http://bgpupdates.potaroo.net/instability/bgpupd.html

The BGP Instability Report

The BGP Instability Report is updated daily. This report was generated on 19 November 2007 02:10 (UTC+1000)

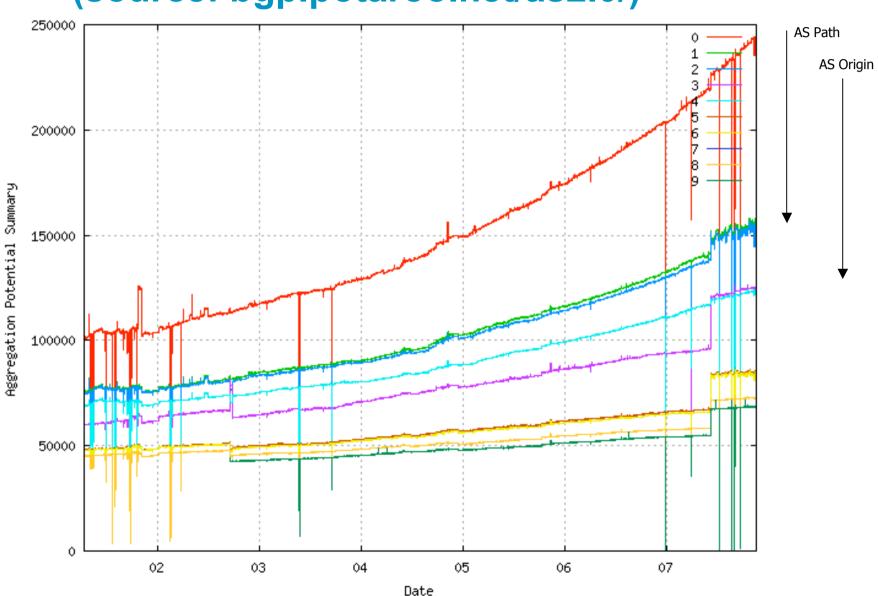
50 Most active ASes for the past 31 days

RANK	ASN	UPDs	%	Prefixes	UPDs/Prefix	AS NAME
1	8452	297064	3.60%	302	983.66	TEDATA TEDATA
2	16637	176716	2.14%	71	2488.96	MTNNS-AS
3	8866	171799	2.08%	281	611.38	BTC-AS Bulgarian Telecommunication Company Plc.
4	9583	109750	1.33%	1158	94.78	SIFY-AS-IN Sify Limited
5	33783	100157	1.21%	128	782.48	EEPAD
6	288	95411	1.16%	121	788.52	European Space Agency
7	5583	85536	1.04%	63	1357.71	GIPNL Equant Benelux AS
8	3215	67691	0.82%	652	103.82	AS3215 France Telecom - Orange
9	9498	62137	0.75%	1071	58.02	BBIL-AP BHARTI BT INTERNET LTD.
10	17540	54919	0.67%	1	54919.00	MTL-AP Modern Terminals Limited
11	4621	54763	0.66%	150	365.09	UNSPECIFIED UNINET-TH
12	14390	53250	0.65%	56	950.89	CORENET - Coretel America, Inc.
13	4755	50962	0.62%	1504	33.88	VSNL-AS Videsh Sanchar Nigam Ltd. Autonomous System
14	8151	48720	0.59%	1319	36.94	Uninet S.A. de C.V.
15	12975	46709	0.57%	80	583.86	PALTEL-AS PALTEL Autonomous System
16	24731	44085	0.53%	47	937.98	ASN-NESMA National Engineering Services and Marketing Company Ltd. (NESMA)
17	26829	43445	0.53%	1	43445.00	YKK-USA - YKK USA,INC
18	702	42445	0.51%	616	68.90	AS702 Verizon Business EMEA - Commercial IP service provider in Europe
19	9835	40758	0.49%	128	318.42	GITS-TH-AS-AP Government Information Technology Services
20	17974	37235	0.45%	535	69.60	TELKOMNET-AS2-AP PT Telekomunikasi Indonesia
21	4750	35626	0.43%	228	156.25	CSLOXINFO-ISP-AS-AP CSLOXINFO Public Company Limited.
22	9829	32905	0.40%	770	42.73	BSNL-NIB National Internet Backbone
23	7545	31349	0.38%	744	42.14	TPG-INTERNET-AP TPG Internet Pty Ltd

50 Most active Prefixes for the past 31 days

RANK	PREFIX	UPDs	%	Origin AS AS NAME
1	192.96.13.0/24	83258	0.95%	16637 MTNNS-AS
2	192.96.14.0/24	83250	0.95%	16637 MTNNS-AS
3	83.228.103.0/24	66715	0.76%	8866 BTC-AS Bulgarian Telecommunication Company Plc.
4	203.83.127.0/24	54919	0.63%	17540 MTL-AP Modern Terminals Limited
5	209.163.125.0/24	52166	0.59%	14390 CORENET - Coretel America, Inc.
6	125.23.208.0/20	46352		9498 BBIL-AP BHARTI BT INTERNET LTD.
7	12.108.254.0/24	43445	0.49%	26829 YKK-USA - YKK USA,INC
8	83.228.61.0/24	32708		8866 BTC-AS Bulgarian Telecommunication Company Plc.
9	83.228.59.0/24	32694		8866 BTC-AS Bulgarian Telecommunication Company Plc.
10	83.228.71.0/24	31792		8866 BTC-AS Bulgarian Telecommunication Company Plc.
11	196.219.236.0/24	23705		8452 TEDATA TEDATA
12	196.219.234.0/24	23694		8452 TEDATA TEDATA
13	196.219.235.0/24	23677		8452 TEDATA TEDATA
14	196.219.244.0/24	23674		8452 TEDATA TEDATA
15	196.219.249.0/24			8452 TEDATA TEDATA
16	196.219.245.0/24			8452 TEDATA TEDATA
17	196.219.248.0/24			8452 TEDATA TEDATA
18	196.219.246.0/24			8452 TEDATA TEDATA
19	196.219.247.0/24			8452 TEDATA TEDATA
	213.158.189.0/24	23544		8452 TEDATA TEDATA
	81.10.120.0/24	22946		8452 TEDATA TEDATA
	210.18.10.0/24	20826		9583 SIFY-AS-IN Sify Limited
	221.135.22.0/24	18137		9583 SIFY-AS-IN Sify Limited
24	221.135.113.0/24			9583 SIFY-AS-IN Sify Limited
25	12.106.30.0/24	15977		22072 - DEFINITYHEALTH - Definity Health
26	194.209.8.0/24	15212		3303 SWISSCOM Swisscom Solutions Ltd
	90.80.16.0/24	14249		43830 ADECCO-ASN Adecco IT Services
28	89.4.131.0/24	12620	0.14%	24731 ASN-NESMA National Engineering Services and Marketing Company Ltd. (NESMA)

Aggregation Potential (source: bgp.potaroo.net/as2.0/)



Aggregation Summary

Aggregation on the Internet could be MUCH better

35% saving on Internet routing table size is quite feasible

Tools are available

Commands on the routers are not hard

CIDR-Report webpage



Receiving Prefixes

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Receiving Prefixes

 There are three scenarios for receiving prefixes from other ASNs

Customer talking BGP

Peer talking BGP

Upstream/Transit talking BGP

 Each has different filtering requirements and need to be considered separately

Receiving Prefixes: From Customers

- ISPs should only accept prefixes which have been assigned or allocated to their downstream customer
- If ISP has assigned address space to its customer, then the customer IS entitled to announce it back to his ISP
- If the ISP has NOT assigned address space to its customer, then:

Check the five RIR databases to see if this address space really has been assigned to the customer

The tool: whois

Receiving Prefixes: From Customers

Example use of whois to check if customer is entitled to announce address space:

```
pfs-pc$ whois -h whois.apnic.net 202.12.29.0
              202.12.29.0 - 202.12.29.255
inetnum:
              APNTC-AP-AU-BNE
netname:
descr:
              APNIC Pty Ltd - Brisbane Offices + Servers
descr:
              Level 1, 33 Park Rd
              PO Box 2131, Milton
descr:
descr:
              Brisbane, QLD.
country:
              ΑU
admin-c:
              HM2.0-AP
                                Portable – means its an assignment
tech-c:
              NO4-AP
                                to the customer, the customer can
mnt-by:
              APNIC-HM
                                announce it to you
              hm-changed@apnic.net 20030108
changed:
              ASSIGNED PORTABLE
status:
              APNIC
source:
```

Receiving Prefixes: From Customers

RIPE

source:

Example use of whois to check if customer is entitled to announce address space:

```
$ whois -h whois.ripe.net 193.128.2.0
              193.128.2.0 - 193.128.2.15
inetnum:
descr:
              Wood Mackenzie
country:
              GB
admin-c:
              DB635-RIPE
                                        ASSIGNED PA - means that it is
tech-c:
              DB635-RIPE
                                        Provider Aggregatable address space
status:
              ASSIGNED PA
                                        and can only be used for connecting
mnt-by:
              AS1849-MNT
                                        to the ISP who assigned it
              davids@uk.uu.net 20020211
changed:
              RIPE
source:
              193.128.0.0/14
route:
descr:
              PIPEX-BLOCK1
origin:
              AS1849
              routing@uk.uu.net
notify:
mnt-by:
              AS1849-MNT
changed:
              beny@uk.uu.net 20020321
```

Receiving Prefixes: From Peers

 A peer is an ISP with whom you agree to exchange prefixes you originate into the Internet routing table

Prefixes you accept from a peer are only those they have indicated they will announce

Prefixes you announce to your peer are only those you have indicated you will announce

Receiving Prefixes: From Peers

Agreeing what each will announce to the other:

Exchange of e-mail documentation as part of the peering agreement, and then ongoing updates

OR

Use of the Internet Routing Registry and configuration tools such as the IRRToolSet

www.isc.org/sw/IRRToolSet/

Receiving Prefixes: From Upstream/Transit Provider

- Upstream/Transit Provider is an ISP who you pay to give you transit to the WHOLE Internet
- Receiving prefixes from them is not desirable unless really necessary

special circumstances – see later

Ask upstream/transit provider to either:

originate a default-route

OR

announce one prefix you can use as default

Receiving Prefixes: From Upstream/Transit Provider

 If necessary to receive prefixes from any provider, care is required

```
don't accept RFC1918 etc prefixes

ftp://ftp.rfc-editor.org/in-notes/rfc3330.txt

don't accept your own prefixes

don't accept default (unless you need it)

don't accept prefixes longer than /24
```

Check Rob Thomas' list of "bogons"

http://www.cymru.com/Documents/bogon-list.html

Receiving Prefixes

Paying attention to prefixes received from customers, peers and transit providers assists with:

The integrity of the local network

The integrity of the Internet

Responsibility of all ISPs to be good Internet citizens



Preparing the network

Before we begin...

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Preparing the Network

- We will deploy BGP across the network before we try and multihome
- BGP will be used therefore an ASN is required
- If multihoming to different ISPs, public ASN needed: Either go to upstream ISP who is a registry member, or Apply to the RIR yourself for a one off assignment, or Ask an ISP who is a registry member, or Join the RIR and get your own IP address allocation too (this option strongly recommended)!

Preparing the Network Initial Assumptions

- The network is not running any BGP at the moment single statically routed connection to upstream ISP
- The network is not running any IGP at all Static default and routes through the network to do "routing"

Preparing the Network First Step: IGP

- Decide on an IGP: OSPF or ISIS ©
- Assign loopback interfaces and /32 address to each router which will run the IGP

Loopback is used for OSPF and BGP router id anchor Used for iBGP and route origination

Deploy IGP (e.g. OSPF)

IGP can be deployed with NO IMPACT on the existing static routing

e.g. OSPF distance might be 110m static distance is 1

Smallest distance wins

Preparing the Network IGP (cont)

Be prudent deploying IGP – keep the Link State Database Lean!

Router loopbacks go in IGP

WAN point to point links go in IGP

(In fact, any link where IGP dynamic routing will be run should go into IGP)

Summarise on area/level boundaries (if possible) – i.e. think about your IGP address plan

Preparing the Network IGP (cont)

Routes which don't go into the IGP include:

Dynamic assignment pools (DSL/Cable/Dial)

Customer point to point link addressing

(using next-hop-self in iBGP ensures that these do NOT need to be in IGP)

Static/Hosting LANs

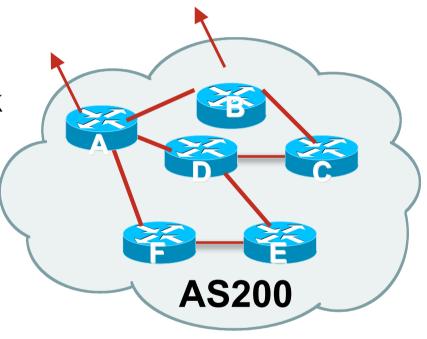
Customer assigned address space

Anything else not listed in the previous slide

Preparing the Network Second Step: iBGP

 Second step is to configure the local network to use iBGP

- iBGP can run on
 all routers, or
 a subset of routers, or
 just on the upstream edge
- iBGP must run on all routers which are in the transit path between external connections

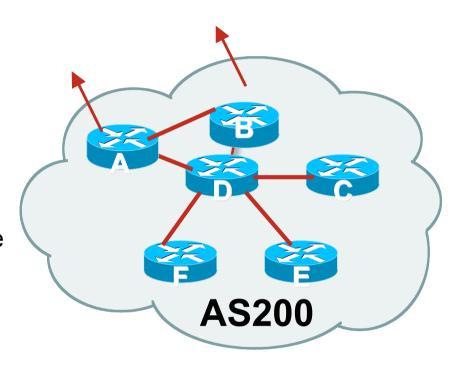


Preparing the Network Second Step: iBGP (Transit Path)

- iBGP must run on all routers which are in the transit path between external connections
- Routers C, E and F are not in the transit path

Static routes or IGP will suffice

 Router D is in the transit path
 Will need to be in iBGP mesh, otherwise routing loops will result



Preparing the Network Layers

Typical SP networks have three layers:

Core – the backbone, usually the transit path

Distribution – the middle, PoP aggregation layer

Aggregation – the edge, the devices connecting customers

Preparing the Network Aggregation Layer

iBGP is optional

Many ISPs run iBGP here, either partial routing (more common) or full routing (less common)

Full routing is not needed unless customers want full table

Partial routing is cheaper/easier, might usually consist of internal prefixes and, optionally, external prefixes to aid external load balancing

Communities and peer-groups make this administratively easy

Many aggregation devices can't run iBGP

Static routes from distribution devices for address pools IGP for best exit

Preparing the Network Distribution Layer

Usually runs iBGP

Partial or full routing (as with aggregation layer)

But does not have to run iBGP

IGP is then used to carry customer prefixes (does not scale)

IGP is used to determine nearest exit

Networks which plan to grow large should deploy iBGP from day one

Migration at a later date is extra work

No extra overhead in deploying iBGP, indeed IGP benefits

Preparing the Network Core Layer

- Core of network is usually the transit path
- iBGP necessary between core devices
 Full routes or partial routes:
 Transit ISPs carry full routes in core
 Edge ISPs carry partial routes only
- Core layer includes AS border routers

Preparing the Network iBGP Implementation

Decide on:

Best iBGP policy

Will it be full routes everywhere, or partial, or some mix?

iBGP scaling technique

Community policy?

Route-reflectors?

Techniques such as peer groups and peer templates?

Preparing the Network iBGP Implementation

Then deploy iBGP:

Step 1: Introduce iBGP mesh on chosen routers make sure that iBGP distance is greater than IGP distance (it usually is)

Step 2: Install "customer" prefixes into iBGP

Check! Does the network still work?

Step 3: Carefully remove the static routing for the prefixes now in IGP and iBGP

Check! Does the network still work?

Step 4: Deployment of eBGP follows

Preparing the Network iBGP Implementation

Install "customer" prefixes into iBGP?

- Customer assigned address space
 Network statement/static route combination
 Use unique community to identify customer assignments
- Customer facing point-to-point links
 - Redistribute connected through filters which only permit point-to-point link addresses to enter iBGP
 - Use a unique community to identify point-to-point link addresses (these are only required for your monitoring system)
- Dynamic assignment pools & local LANs
 - Simple network statement will do this
 - Use unique community to identify these networks

Preparing the Network iBGP Implementation

Carefully remove static routes?

Work on one router at a time:

Check that static route for a particular destination is also learned by the iBGP

If so, remove it

If not, establish why and fix the problem

(Remember to look in the RIB, not the FIB!)

- Then the next router, until the whole PoP is done
- Then the next PoP, and so on until the network is now dependent on the IGP and iBGP you have deployed

Preparing the Network Completion

Previous steps are NOT flag day steps

Each can be carried out during different maintenance periods, for example:

Step One on Week One

Step Two on Week Two

Step Three on Week Three

And so on

And with proper planning will have NO customer visible impact at all

Preparing the Network Example Two

- The network is not running any BGP at the moment single statically routed connection to upstream ISP
- The network is running an IGP though All internal routing information is in the IGP By IGP, OSPF or ISIS is assumed

Preparing the Network IGP

- If not already done, assign loopback interfaces and /32 addresses to each router which is running the IGP
 - Loopback is used for OSPF and BGP router id anchor
 - Used for iBGP and route origination
- Ensure that the loopback /32s are appearing in the IGP

Preparing the Network iBGP

- Go through the iBGP decision process as in Example One
- Decide full or partial, and the extent of the iBGP reach in the network

Preparing the Network iBGP Implementation

- Then deploy iBGP:
 - Step 1: Introduce iBGP mesh on chosen routers

 make sure that iBGP distance is greater than IGP distance (it usually is)
 - Step 2: Install "customer" prefixes into iBGP
 - Check! Does the network still work?
 - Step 3: Reduce BGP distance to be less than the IGP (so that iBGP routes take priority)
 - Step 4: Carefully remove the "customer" prefixes from the IGP Check! Does the network still work?
 - Step 5: Restore BGP distance to less than IGP
 - Step 6: Deployment of eBGP follows

Preparing the Network iBGP implementation

Install "customer" prefixes into iBGP?

- Customer assigned address space
 - Network statement/static route combination
 - Use unique community to identify customer assignments
- Customer facing point-to-point links
 - Redistribute connected through filters which only permit point-to-point link addresses to enter iBGP
 - Use a unique community to identify point-to-point link addresses (these are only required for your monitoring system)
- Dynamic assignment pools & local LANs
 - Simple network statement will do this
 - Use unique community to identify these networks

Preparing the Network iBGP implementation

Carefully remove "customer" routes from IGP?

Work on one router at a time:

Check that IGP route for a particular destination is also learned by iBGP

If so, remove it from the IGP

If not, establish why and fix the problem

(Remember to look in the RIB, not the FIB!)

- Then the next router, until the whole PoP is done
- Then the next PoP, and so on until the network is now dependent on the iBGP you have deployed

Preparing the Network Completion

Previous steps are NOT flag day steps

Each can be carried out during different maintenance periods, for example:

Step One on Week One

Step Two on Week Two

Step Three on Week Three

And so on

And with proper planning will have NO customer visible impact at all

Preparing the Network Configuration Summary

- IGP essential networks are in IGP
- Customer networks are now in iBGP iBGP deployed over the backbone Full or Partial or Upstream Edge only
- BGP distance is greater than any IGP
- Now ready to deploy eBGP



Configuration Tips

Of passwords, tricks and templates

iBGP and IGPs Reminder!

- Make sure loopback is configured on router iBGP between loopbacks, NOT real interfaces
- Make sure IGP carries loopback /32 address
- Consider the DMZ nets:
 - Use unnumbered interfaces?
 - Use next-hop-self on iBGP neighbours
 - Or carry the DMZ /30s in the iBGP
 - Basically keep the DMZ nets out of the IGP!

iBGP: Next-hop-self

- BGP speaker announces external network to iBGP peers using router's local address (loopback) as nexthop
- Used by many ISPs on edge routers

Preferable to carrying DMZ /30 addresses in the IGP

Reduces size of IGP to just core infrastructure

Alternative to using unnumbered interfaces

Helps scale network

Many ISPs consider this "best practice"

Limiting AS Path Length

 Some BGP implementations have problems with long AS_PATHS

Memory corruption

Memory fragmentation

 Even using AS_PATH prepends, it is not normal to see more than 20 ASes in a typical AS_PATH in the Internet today

The Internet is around 5 ASes deep on average

Largest AS_PATH is usually 16-20 ASNs

Limiting AS Path Length

Some announcements have ridiculous lengths of ASpaths:

This example is an error in one IPv6 implementation

```
*> 194.146.180.0/22 2497 3257 29686 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327 16327
```

This example shows 20 prepends (for no obvious reason)

 If your implementation supports it, consider limiting the maximum AS-path length you will accept

BGP TTL "hack"

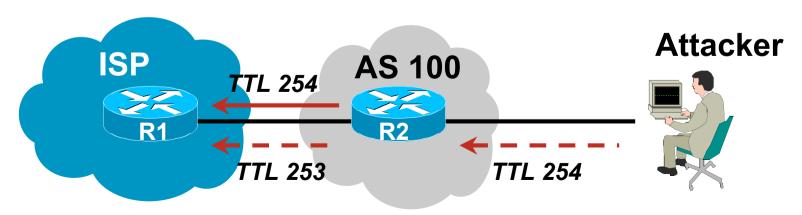
Implement RFC5082 on BGP peerings

(Generalised TTL Security Mechanism)

Neighbour sets TTL to 255

Local router expects TTL of incoming BGP packets to be 254

No one apart from directly attached devices can send BGP packets which arrive with TTL of 254, so any possible attack by a remote miscreant is dropped due to TTL mismatch



BGP TTL "hack"

TTL Hack:

Both neighbours must agree to use the feature TTL check is much easier to perform than MD5 (Called BTSH – BGP TTL Security Hack)

Provides "security" for BGP sessions

In addition to packet filters of course

MD5 should still be used for messages which slip through the TTL hack

See www.nanog.org/mtg-0302/hack.html for more details

Templates

Good practice to configure templates for everything

Vendor defaults tend not to be optimal or even very useful for ISPs

ISPs create their own defaults by using configuration templates

eBGP and iBGP examples follow

Also see Project Cymru's BGP templates

www.cymru.com/Documents

iBGP Template Example

- iBGP between loopbacks!
- Next-hop-self
 Keep DMZ and external point-to-point out of IGP
- Always send communities in iBGP
 Otherwise accidents will happen
- Hardwire BGP to version 4
 Yes, this is being paranoid!

iBGP Template Example continued

Use passwords on iBGP session

Not being paranoid, VERY necessary

It's a secret shared between you and your peer

If arriving packets don't have the correct MD5 hash, they are ignored

Helps defeat miscreants who wish to attack BGP sessions

 Powerful preventative tool, especially when combined with filters and the TTL "hack"

eBGP Template Example

BGP damping

Do **NOT** use it unless you understand the impact Do **NOT** use the vendor defaults without thinking

- Remove private ASes from announcements
 Common omission today
- Use extensive filters, with "backup"
 - Use as-path filters to backup prefix filters
 - Keep policy language for implementing policy, rather than basic filtering
- Use password agreed between you and peer on eBGP session

eBGP Template Example continued

- Use maximum-prefix tracking
 - Router will warn you if there are sudden increases in BGP table size, bringing down eBGP if desired
- Limit maximum as-path length inbound
- Log changes of neighbour state
 - ...and monitor those logs!
- Make BGP admin distance higher than that of any IGP Otherwise prefixes heard from outside your network could override your IGP!!

Summary

- Use configuration templates
- Standardise the configuration
- Be aware of standard "tricks" to avoid compromise of the BGP session
- Anything to make your life easier, network less prone to errors, network more likely to scale
- It's all about scaling if your network won't scale, then it won't be successful



BGP Techniques for Internet Service Providers

Philip Smith <pfs@cisco.com>

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